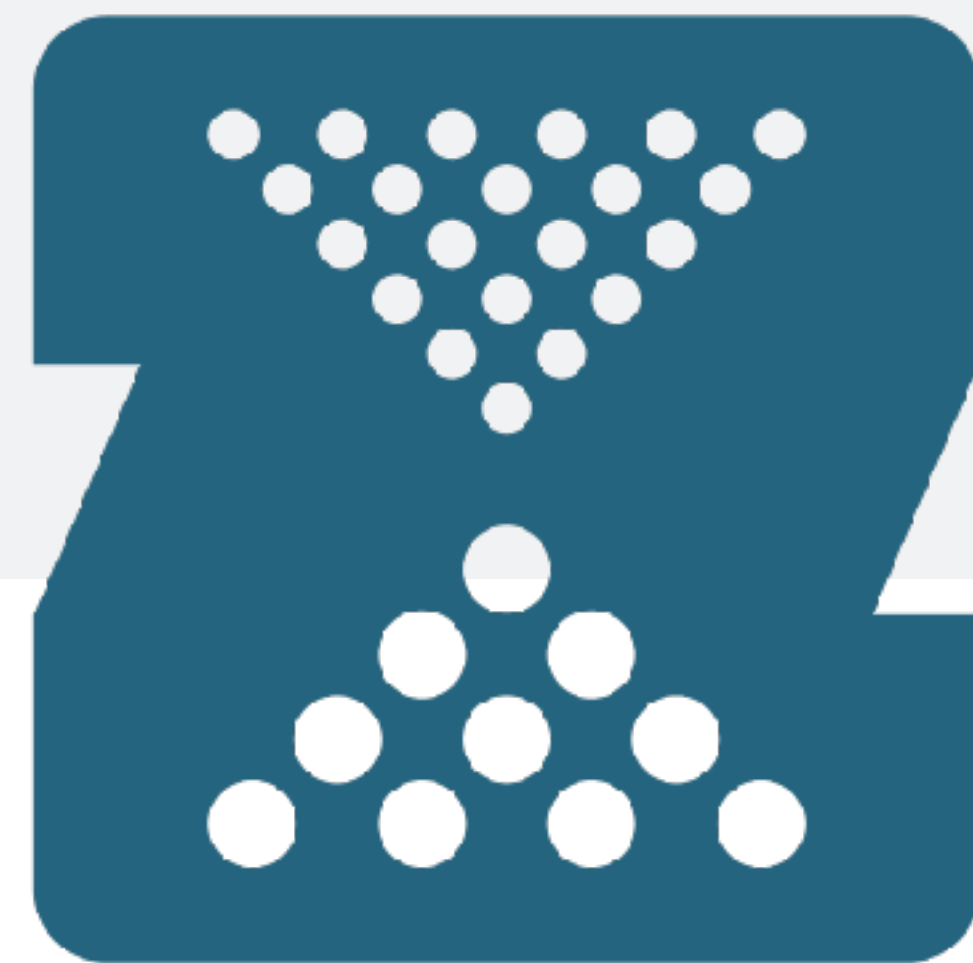
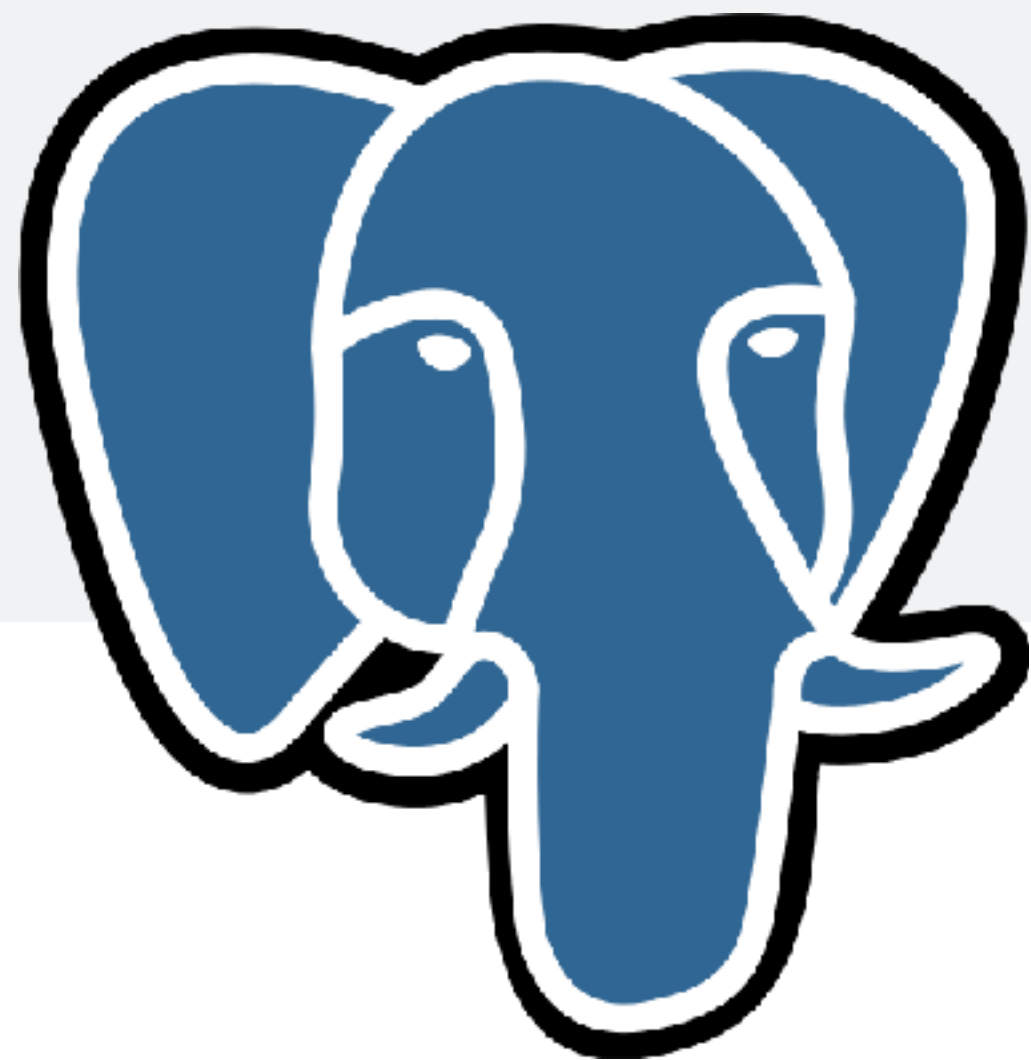
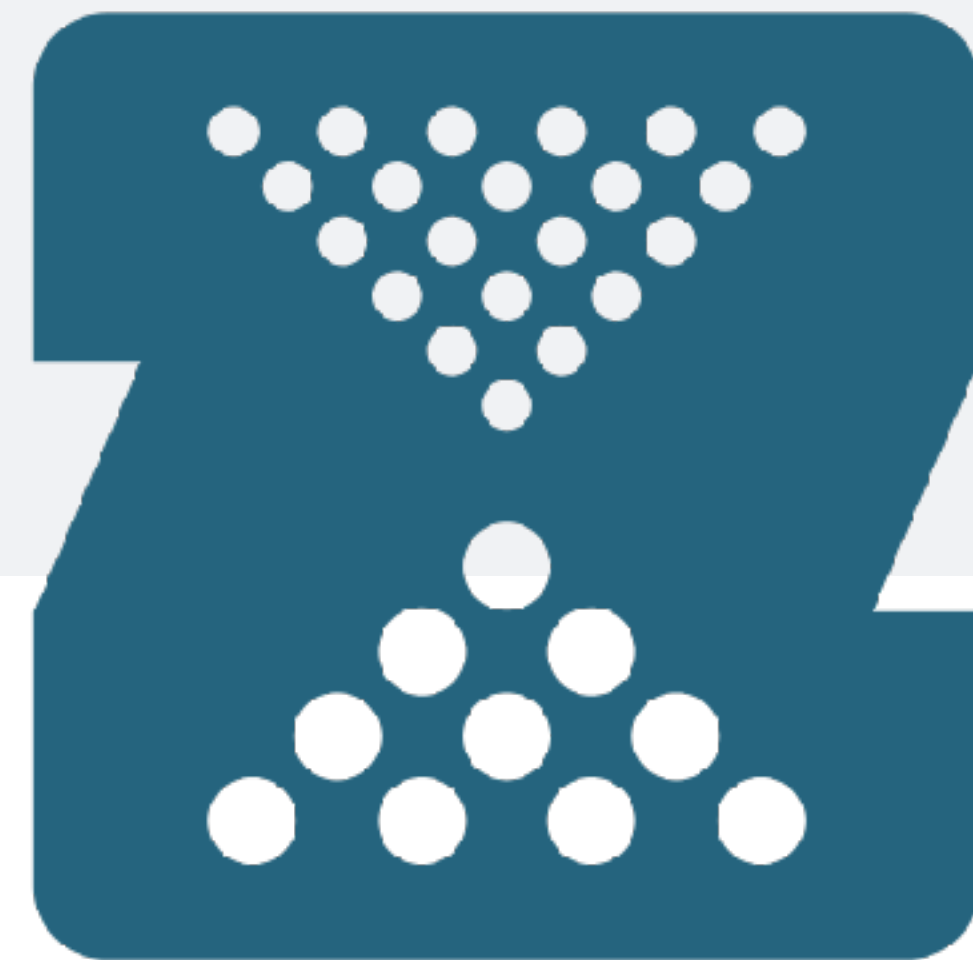
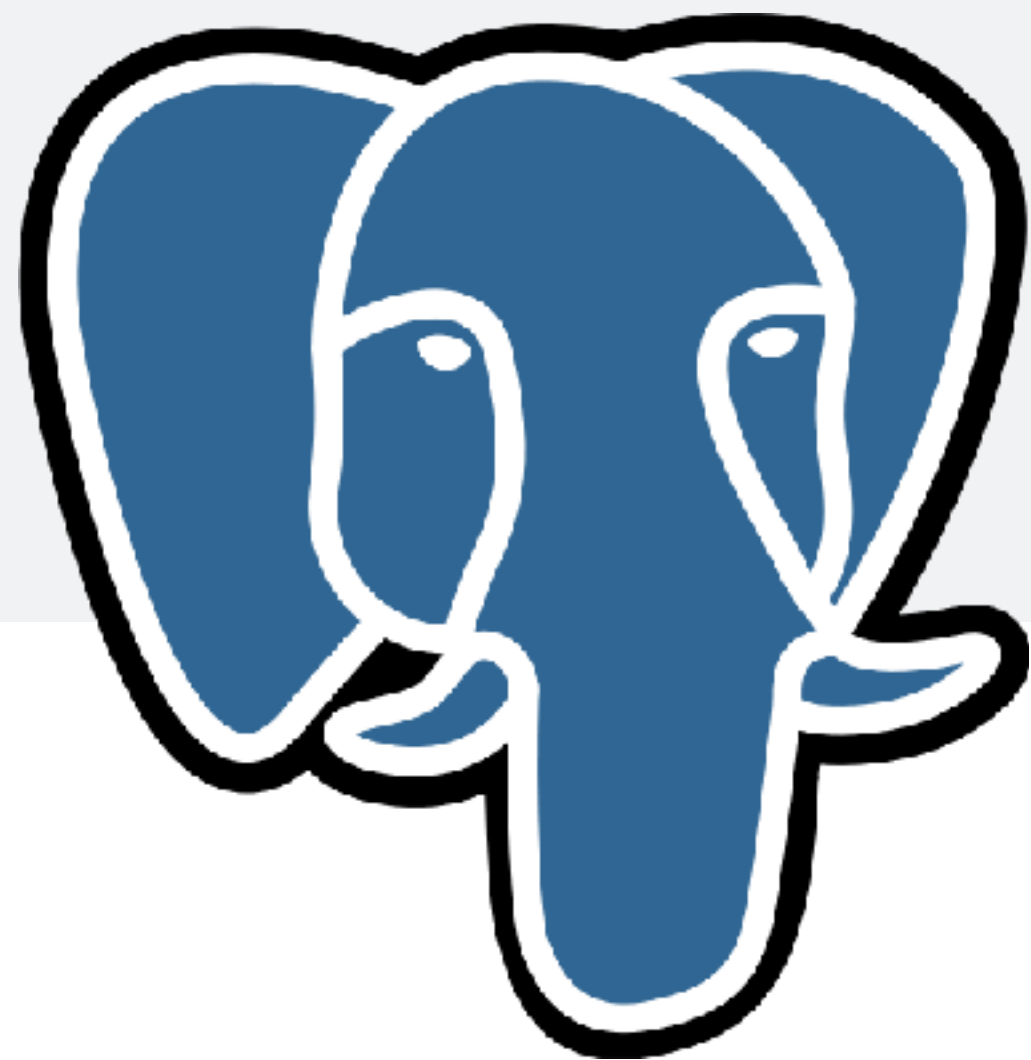


PostgreSQL + ZFS

Best Practices and Standard Procedures



"If you are not using ZFS,
you are losing data*."



Clark's Three Laws

1. When a distinguished but elderly scientist states that something is possible, he is almost certainly right. When he states that something is impossible, he is very probably wrong.
2. The only way of discovering the limits of the possible is to venture a little way past them into the impossible.
3. Any sufficiently advanced technology is indistinguishable from magic.

ZFS is not magic, but it is an incredibly impressive piece of software.

PostgreSQL and ZFS

- Many bits
- Lots of bits
- Huge bits
- It's gunna be great
- Very excited
- We have the best filesystems
- People tell me this is true
- Except the fake media, they didn't tell me this

PostgreSQL and ZFS: It's about the bits and storage, stupid.

- ~~Many bits~~
- ~~Lots of bits~~
- ~~Huge bits~~
- ~~It's gunna be great~~
- ~~Very excited~~
- ~~We have the best filesystems~~
- ~~People tell me this is true~~
- ~~Except the fake media, they didn't tell me this~~

Too soon?

PostgreSQL and ZFS

1. Review PostgreSQL from a storage administrator's perspective
2. Learn what it takes to become a PostgreSQL "backup expert"
3. Dive through a naive block-based filesystem
4. Walk through the a high-level abstraction of ZFS
5. See some examples of how to use ZFS with PostgreSQL
 - Tips
 - Tunables
 - Anecdotes

Some FS minutiae may have been harmed in the making of this talk.
Nit-pick as necessary (preferably after).

PostgreSQL - A Storage Administrator's View

- User-land page cache maintained by PostgreSQL in shared memory
- 8K page size
- Each PostgreSQL table is backed by one or more files in `$PGDATA/`
- Tables larger than 1GB are automatically shared into individual 1GB files
- `pwrite(2)`'s to tables are:
 - append-only if no free pages in the table are available
 - in-place updated if free pages are available in the free-space map
- `pwrite(2)`'s are page-aligned
- Makes heavy use of a Write Ahead Log (WAL), aka an Intent Log

Storage Administration: WAL on Disk

- WAL files are written to sequentially
- append-only IO
- Still 8K page-aligned writes via `pwrite(2)`
- WAL logs are 16MB each, pre-allocated
- WAL logs are never `unlink(2)`'ed, only recycled via `rename(2)`
- Low-latency `pwrite(2)`'s and `fsync(2)` for WAL files is required for good write performance

PostgreSQL - Backups

Traditionally, only two SQL commands that you must know:

1. `pg_start_backup('my_backup')`
2. `${some_random_backup_utility} $PGDATA/`
3. `pg_stop_backup()`

Wait for `pg_start_backup()` to return before backing up `$PGDATA/` directory.

PostgreSQL - Backups

Only ~~two~~^wthree SQL commands that you must know:

1. `CHECKPOINT`
2. `pg_start_backup('my_backup')`
3. `${some_random_backup_utility} $PGDATA/`
4. `pg_stop_backup()`

Manual `CHECKPOINT` if you can't twiddle the args to `pg_start_backup()`.

PostgreSQL - Backups

Only ~~two~~^{three} commands that you must know:

1. ~~CHECKPOINT~~

2. `pg_start_backup('my_backup', true)`

3. `${some_random_backup_utility} $PGDATA/`

4. `pg_stop_backup()`

`pg_start_backup('my_backup', true)`

a.k.a. aggressive checkpointing (vs default perf hit of:

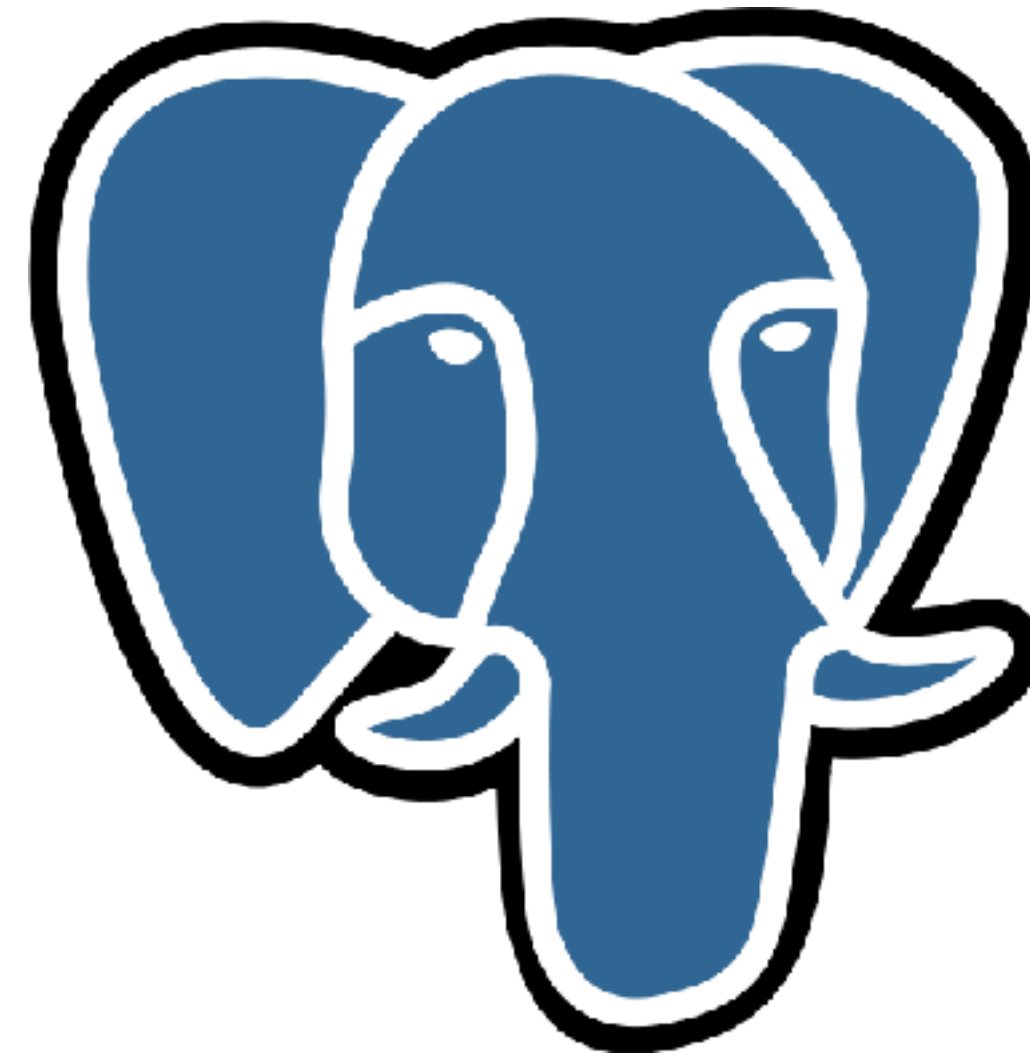
`0.5 * checkpoint_completion_target`)



Achievement unlocked
PostgreSQL Backup Expert



Achievement Pending
ZFS Ninja



TIP: Look for parallels.

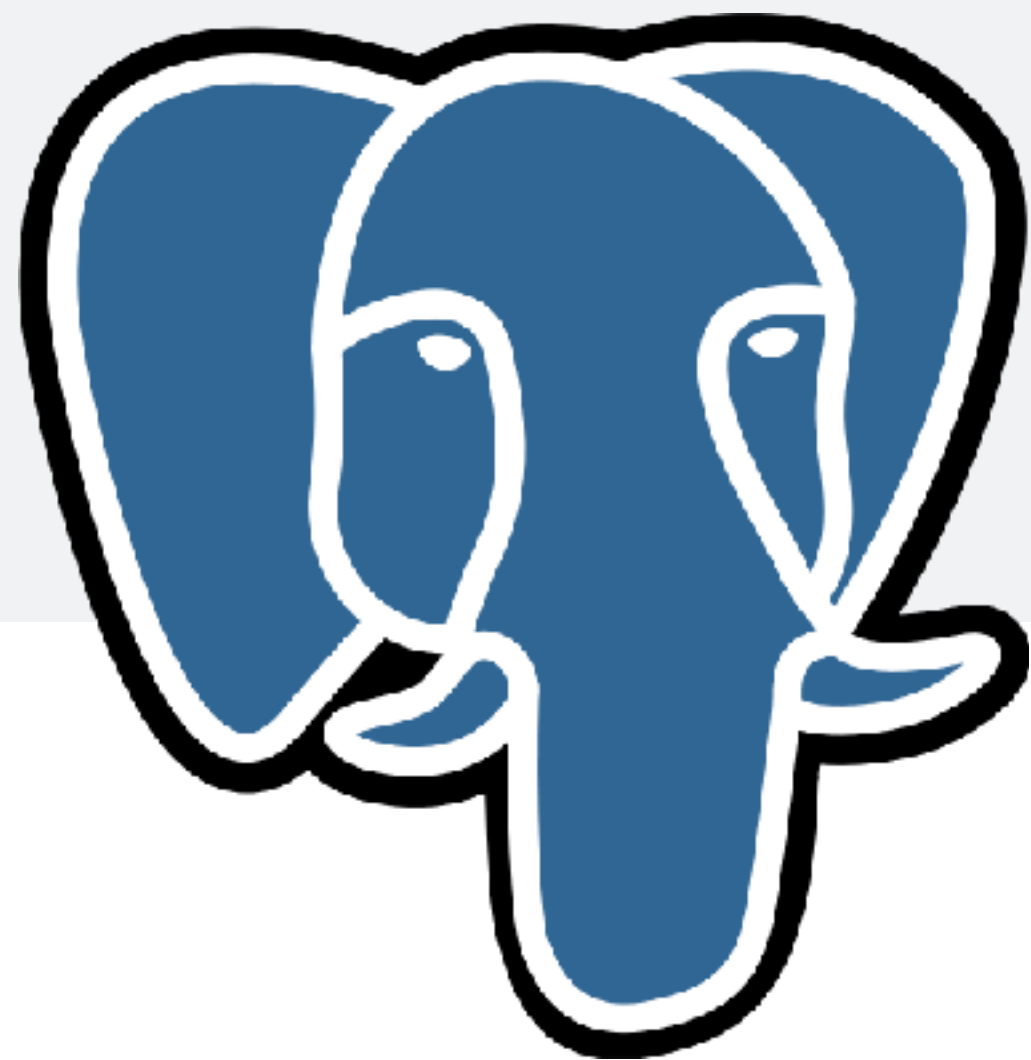
Quick ZFS Primer: Features (read: why you must use ZFS)

- Never inconsistent (no `fsck(8)`'s required, ever)
- Filesystem atomically moves from one consistent state to another consistent state
- All blocks are checksummed
- Compression builtin
- Snapshots are free and unlimited
- Clones are easy
- Changes accumulate in memory, flushed to disk in a transaction
- Redundant metadata (and optionally data)
- Filesystem management independent of physical storage management
- Log-Structured Filesystem
- Copy on Write (COW)

Feature Consequences (read: how your butt gets saved)

- bitrot detected and automatically corrected if possible
 - phantom writes
 - misdirected reads or writes by the drive heads
 - DMA parity errors
 - firmware or driver bugs
 - RAM capacitors aren't refreshed fast enough or with enough power
- Phenomenal sequential and random IO write performance
- Performance increase for sequential reads
- Cost of ownership goes down
- New tricks and tools to solve "data gravity" problems

ELI5: Block Filesystems vs Log Structured Filesystems



Block Filesystems: Top-Down

Userland Application

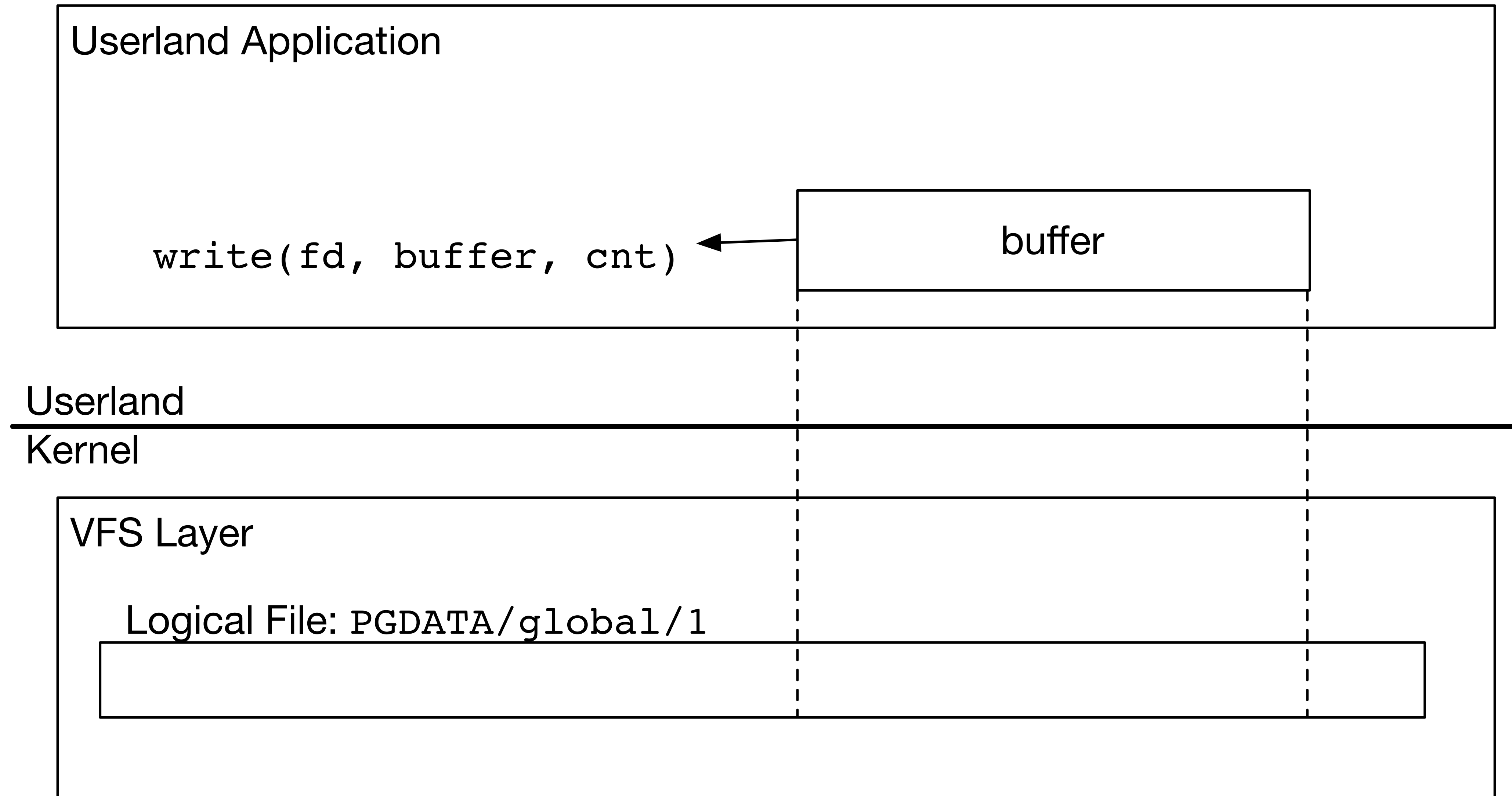
```
write(fd, buffer, cnt)
```

buffer

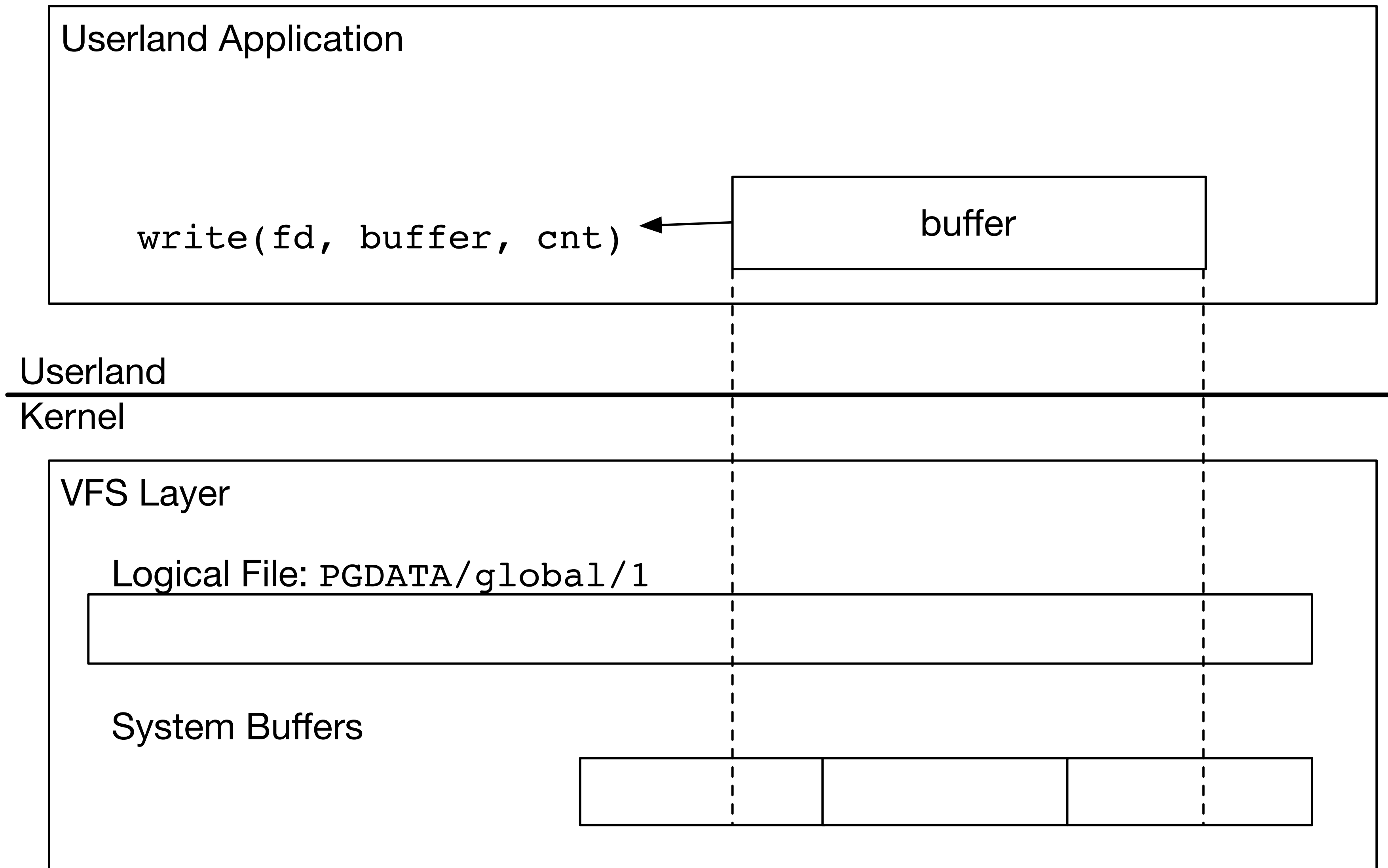
A rectangular box labeled 'buffer' is positioned to the right of the 'write' function call. An arrow points from the right side of this box to the 'buffer' argument in the function call.

Userland

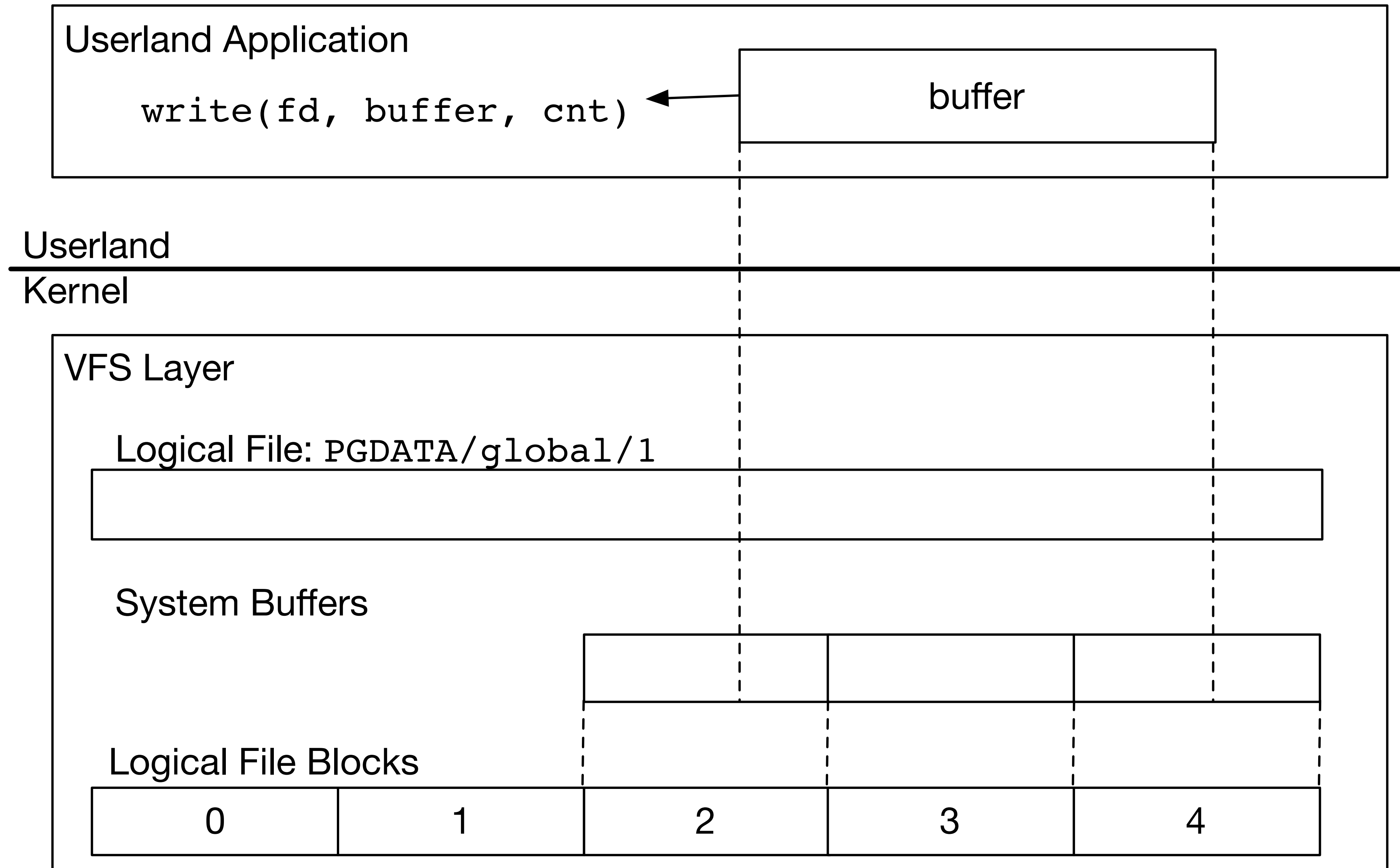
Block Filesystems: Top-Down



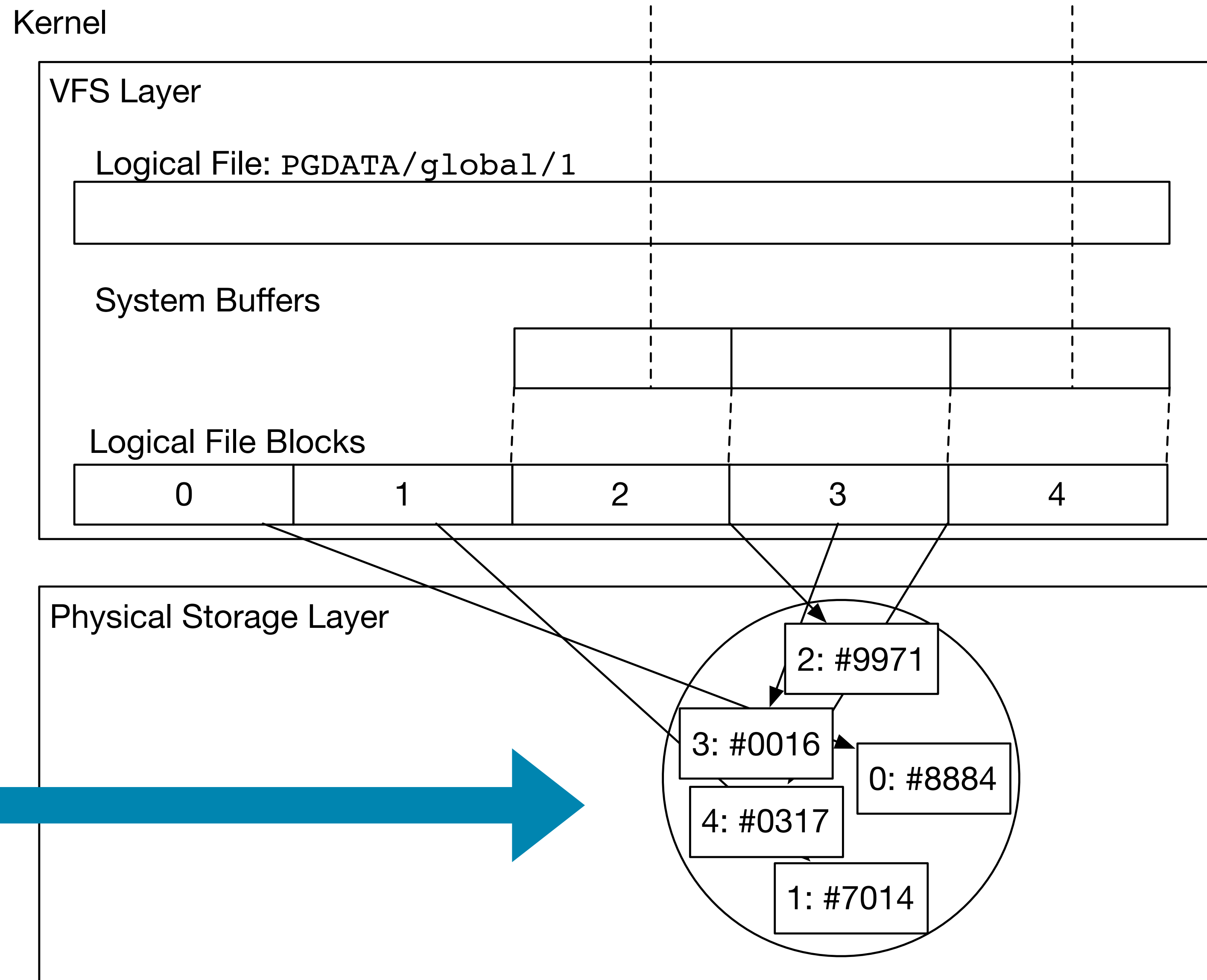
Block Filesystems: Top-Down



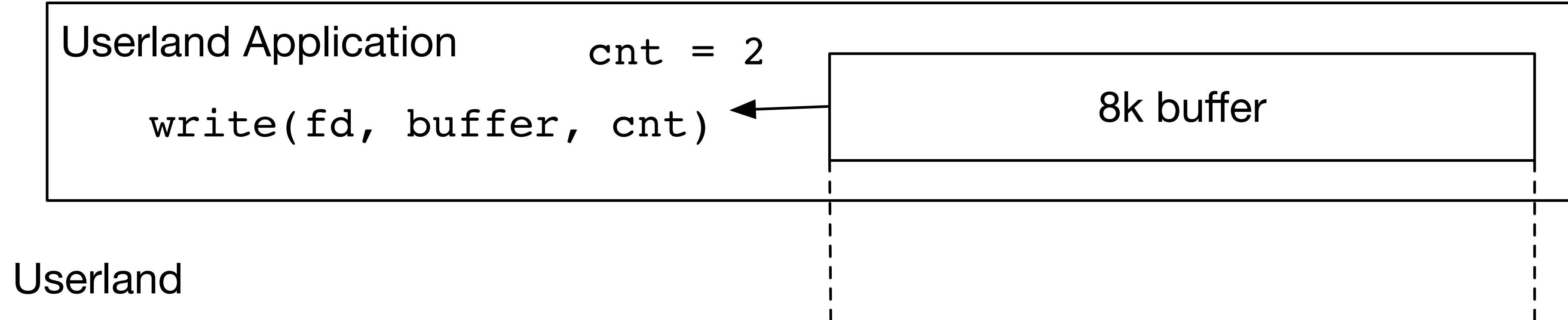
Block Filesystems: Top-Down



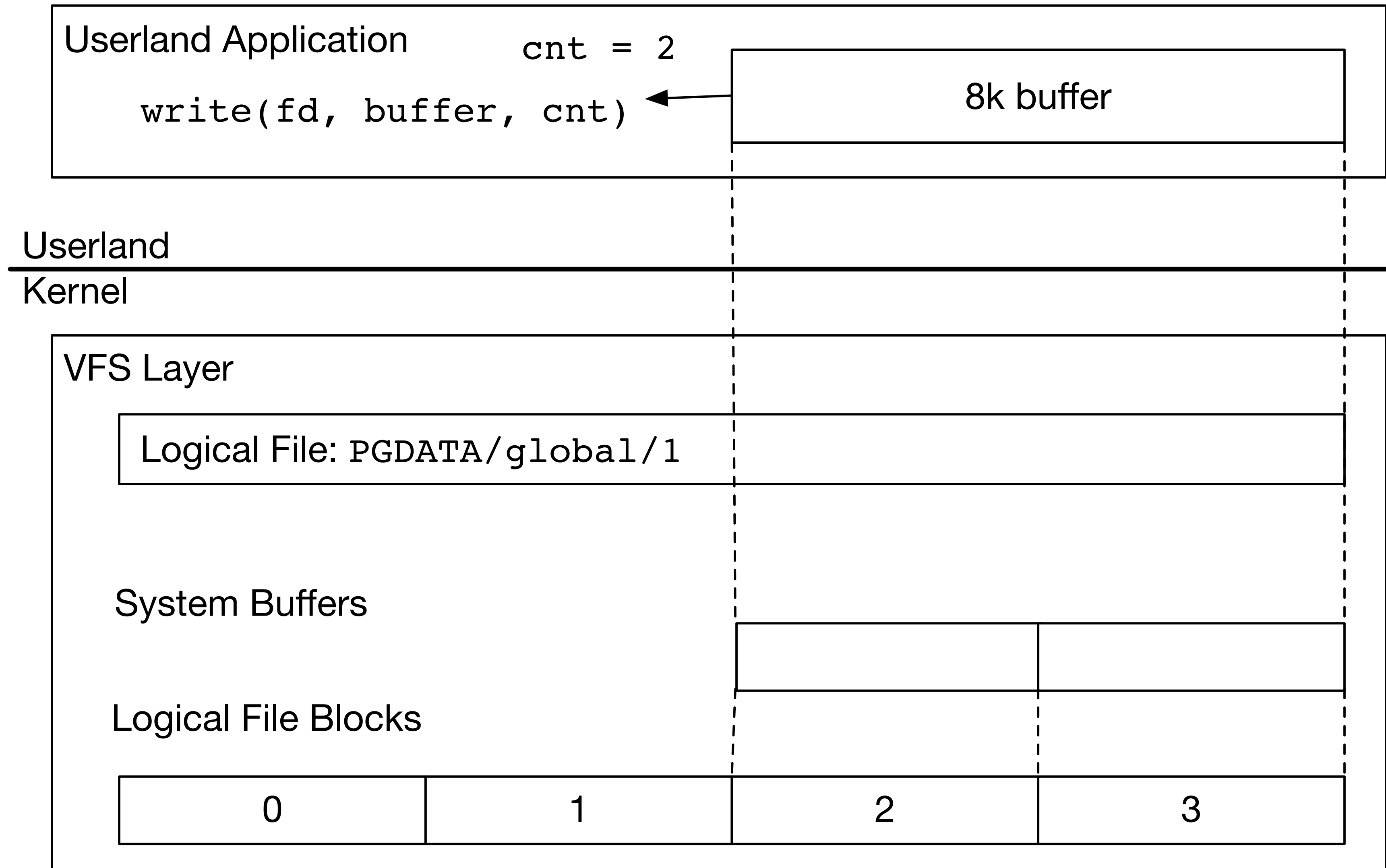
Block Filesystems: Top-Down



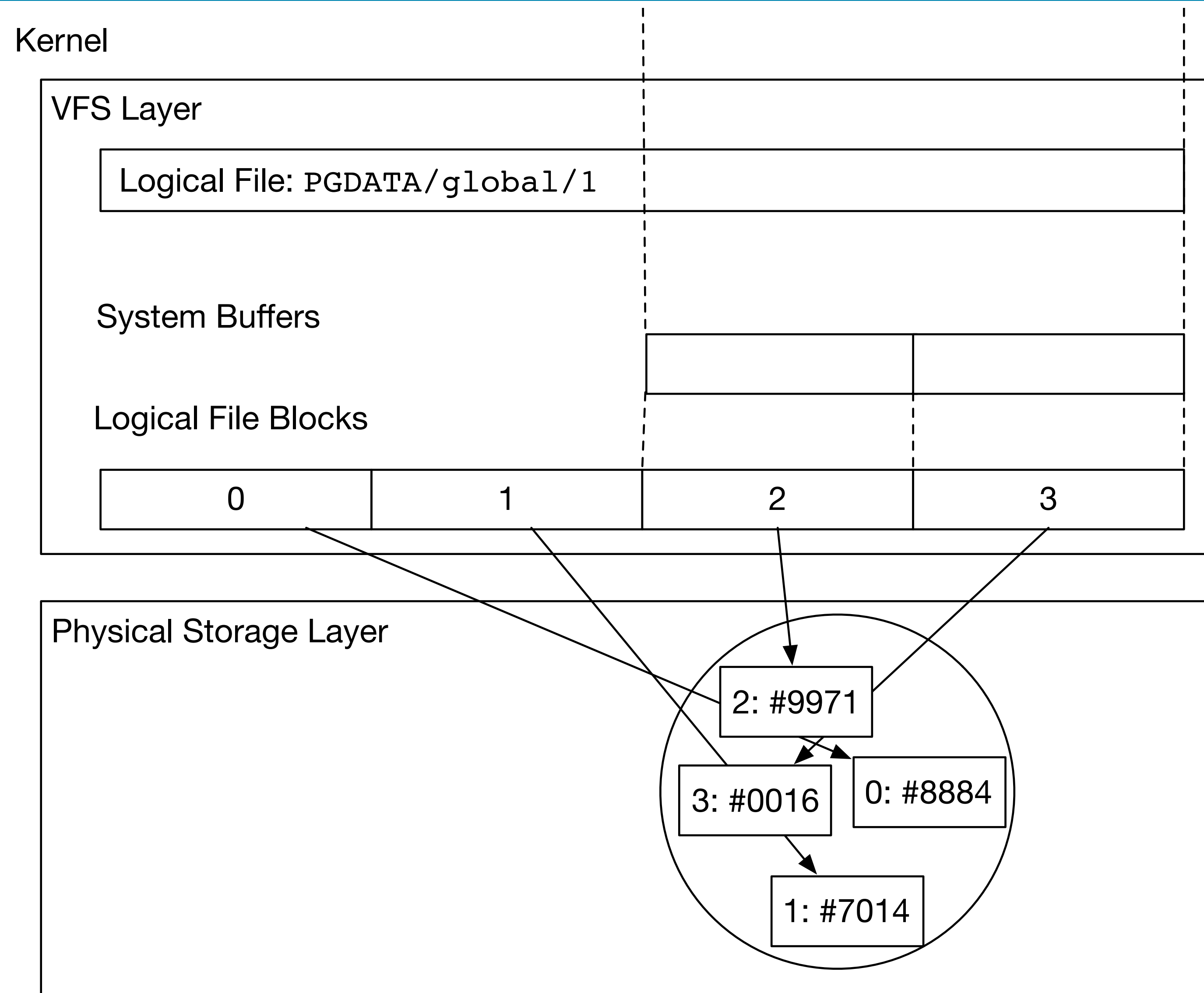
Block Filesystems: PostgreSQL Edition



Block Filesystems: PostgreSQL Edition



Block Filesystems: PostgreSQL Edition

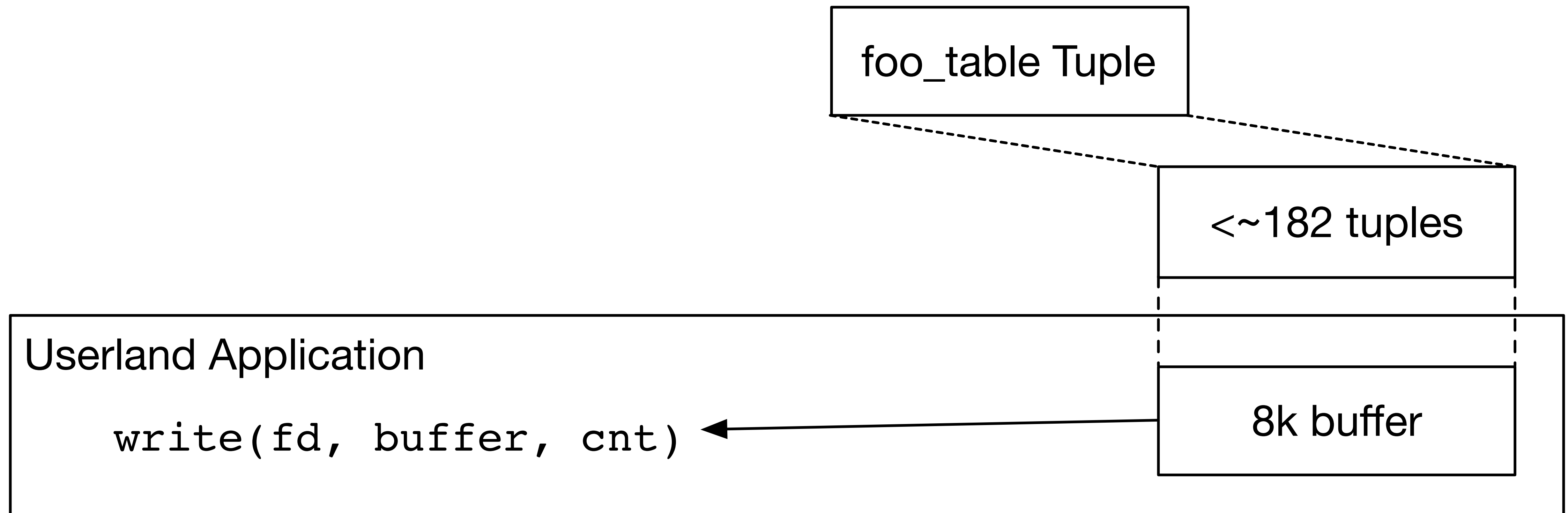


What happens when you twiddle a bool in a row?

```
UPDATE foo_table SET enabled = FALSE WHERE id = 123;
```

Quiz Answer: Write Amplification

```
UPDATE foo_table SET enabled = FALSE WHERE id = 123;
```



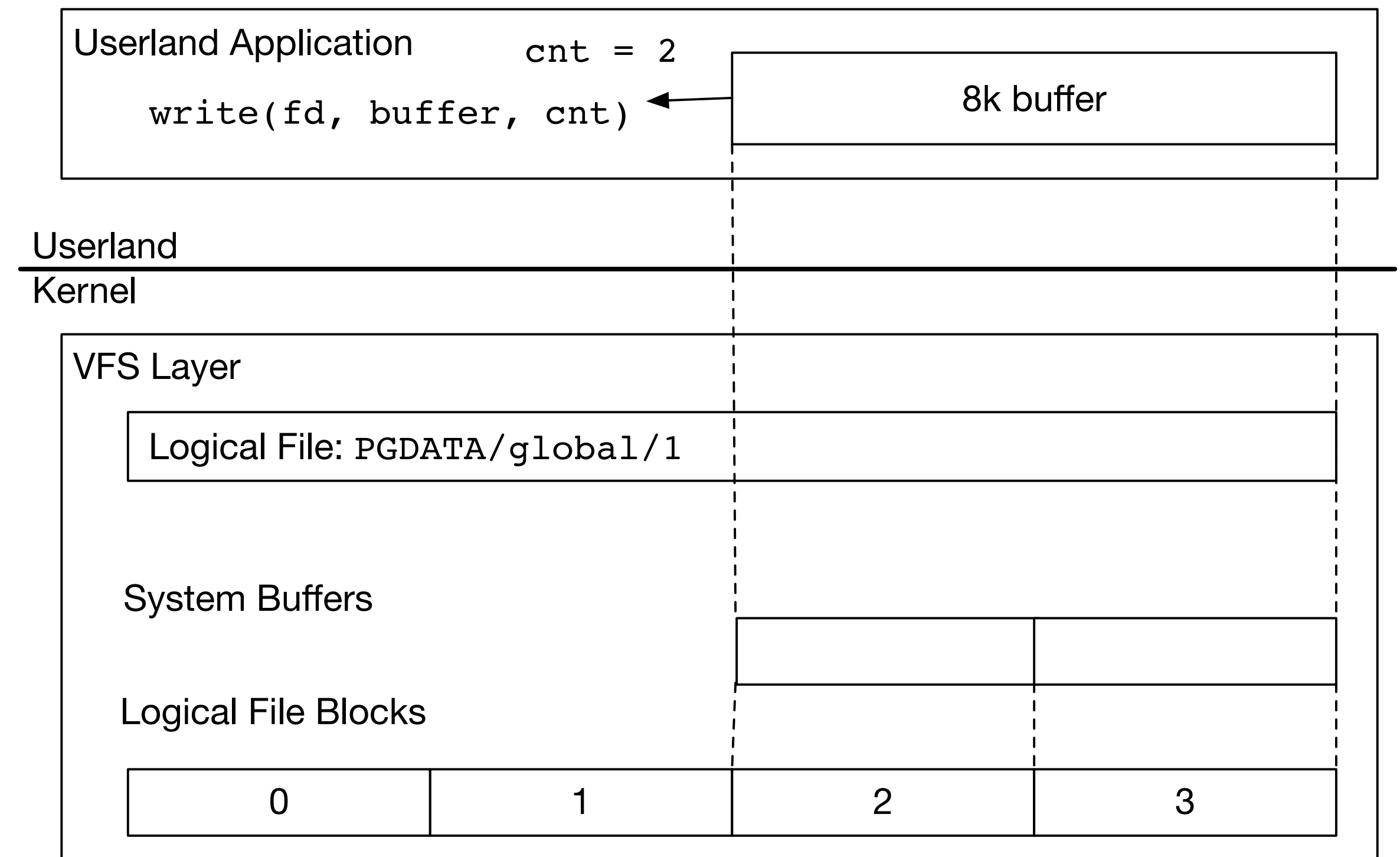
ZFS Tip: postgresql.conf: full_page_writes=off

```
ALTER SYSTEM SET full_page_writes=off;  
CHECKPOINT;  
-- Restart PostgreSQL
```

IMPORTANT NOTE: **full_page_writes=off** interferes with cascading replication

Block Filesystems: PostgreSQL Edition

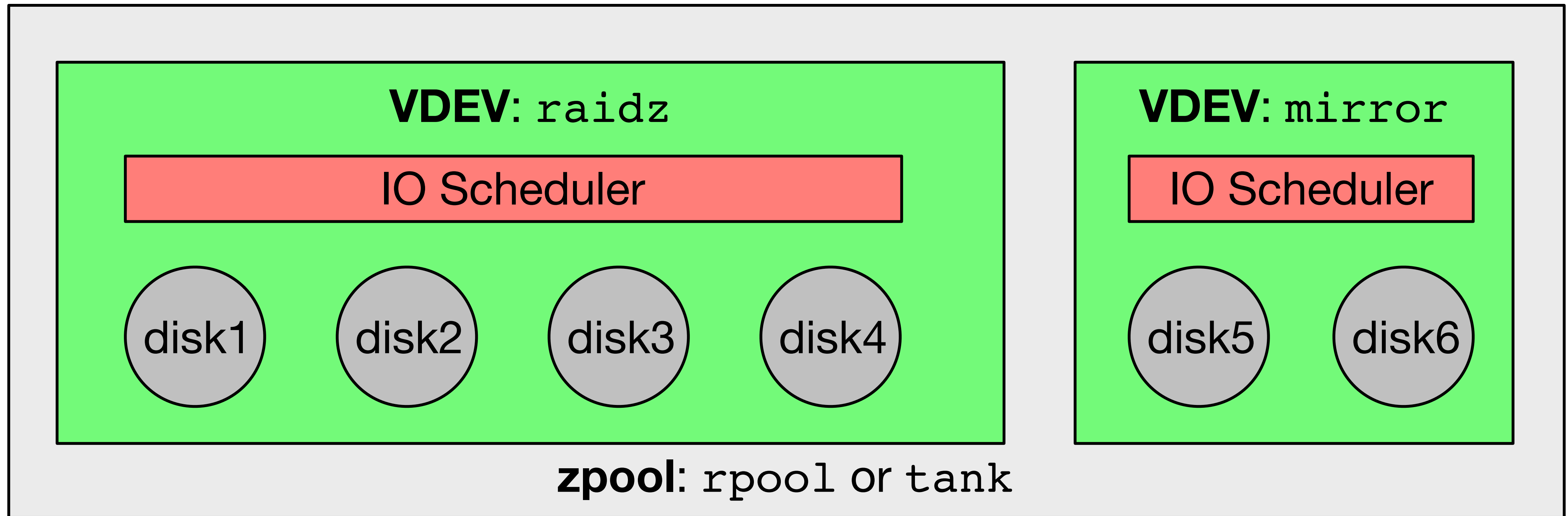
- buffers can be 4K
- disk sectors are 512B - 4K
- ordering of writes is important
- consistency requires complete cooperation and coordination



**Physical Storage is
decoupled
from
Filesystems.**

If you remember one thing from this section,
this is it.

VDEVs On the Bottom

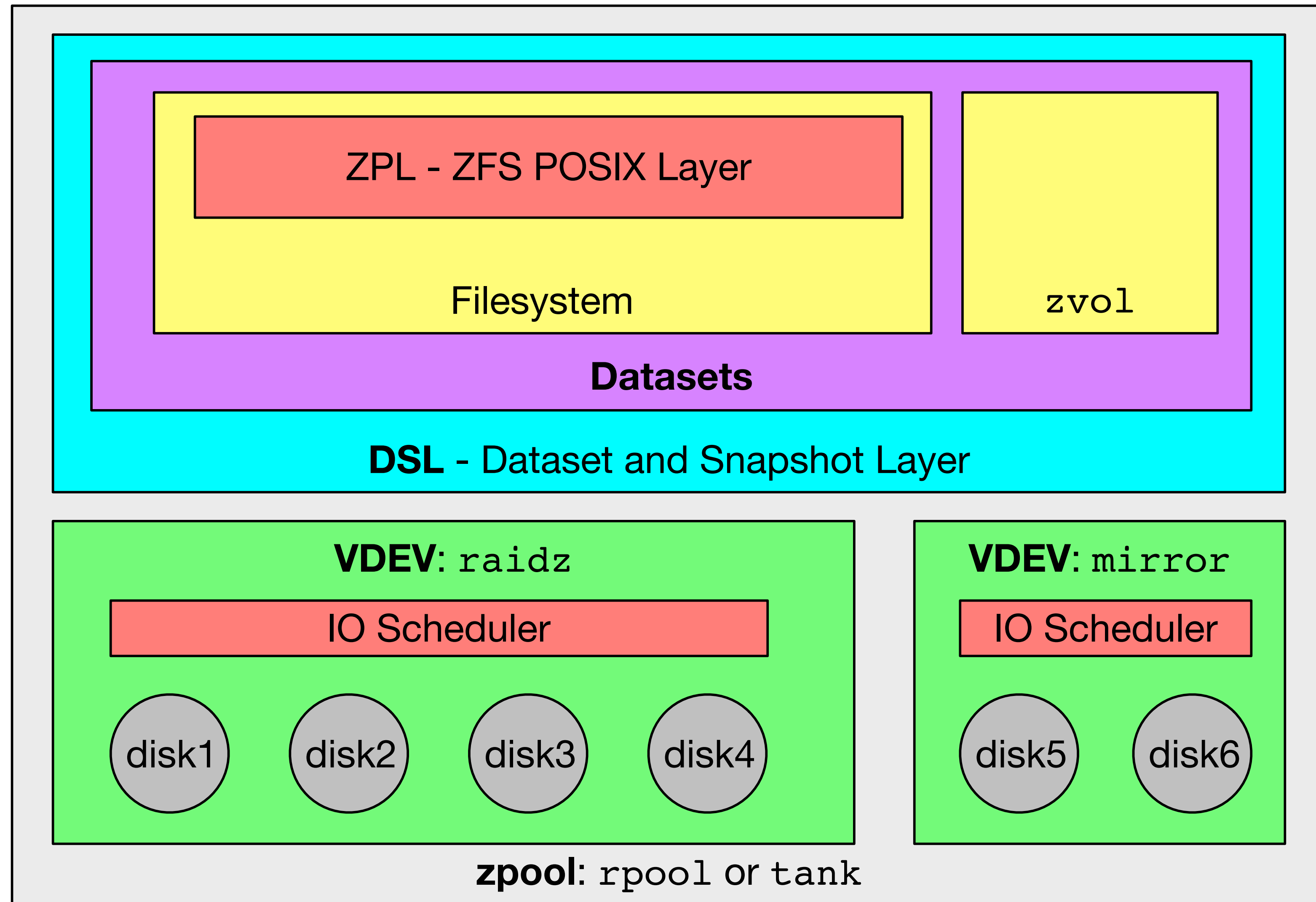


Dataset Name	VFS Mountpoint
tank/ROOT	/
tank/db	/db
tank/ROOT/usr	/usr
tank/local	none
tank/local/etc	/usr/local/etc

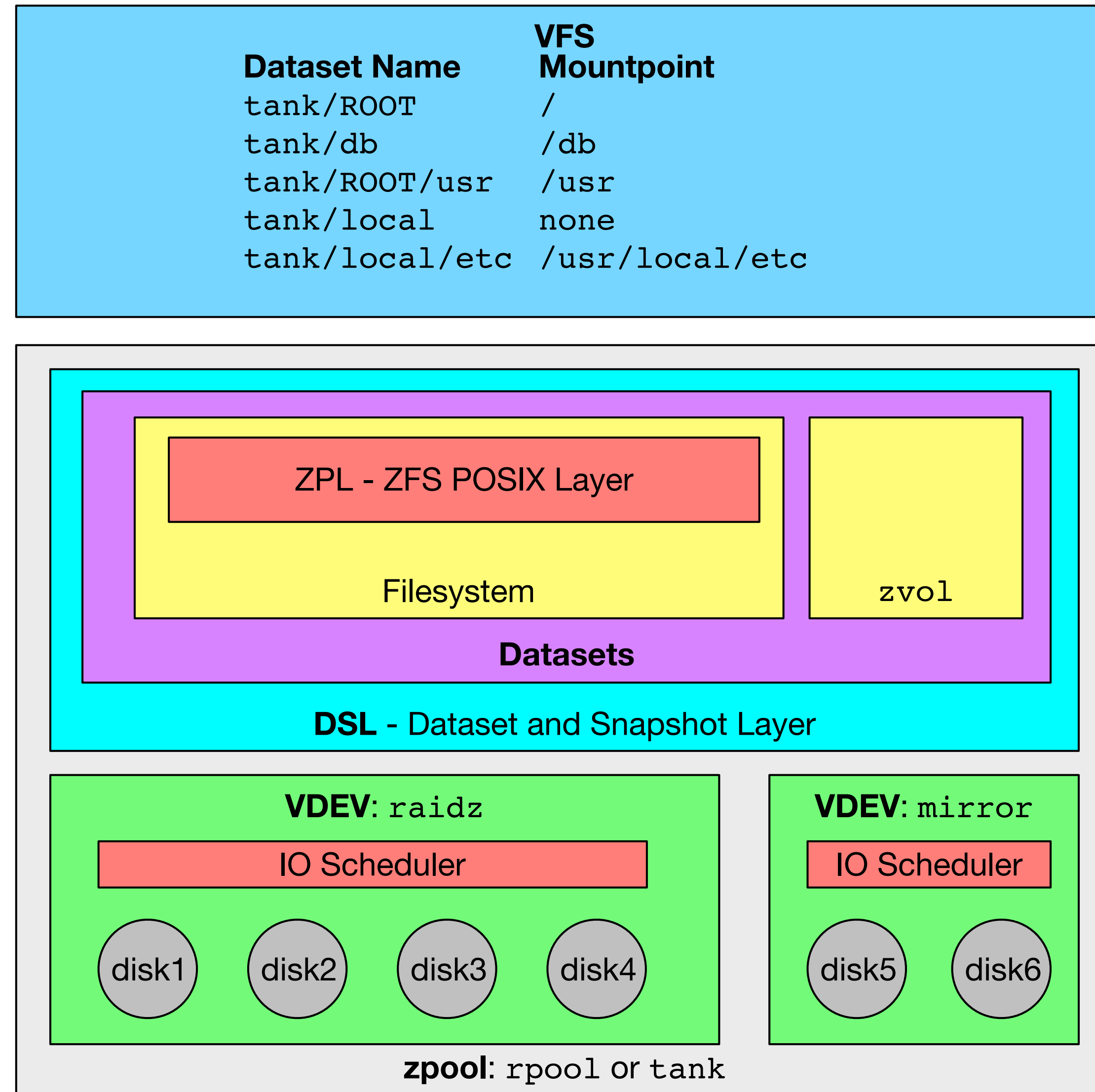
canmount=off



Offensively Over Simplified Architecture Diagram



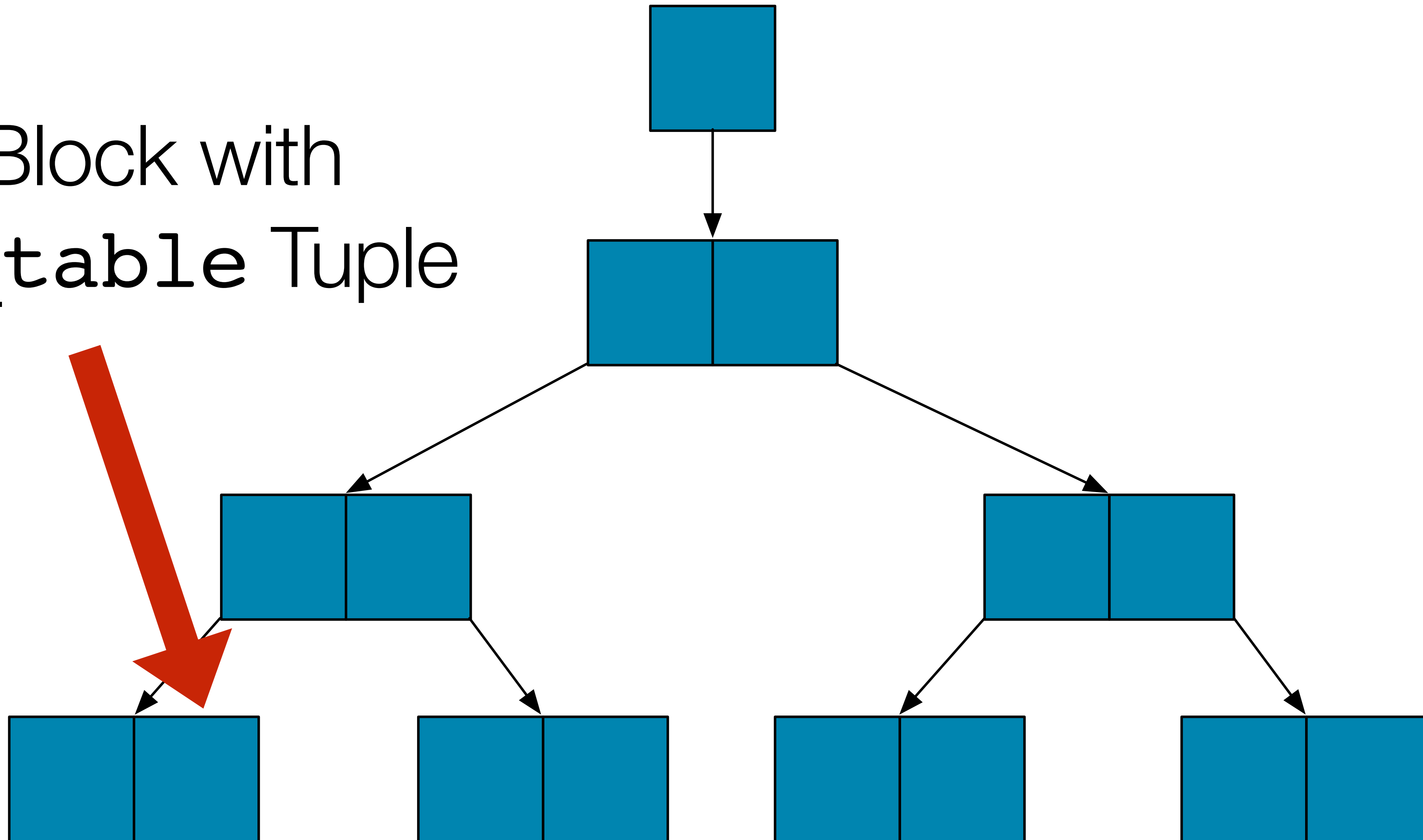
ZFS is magic until you know how it fits together



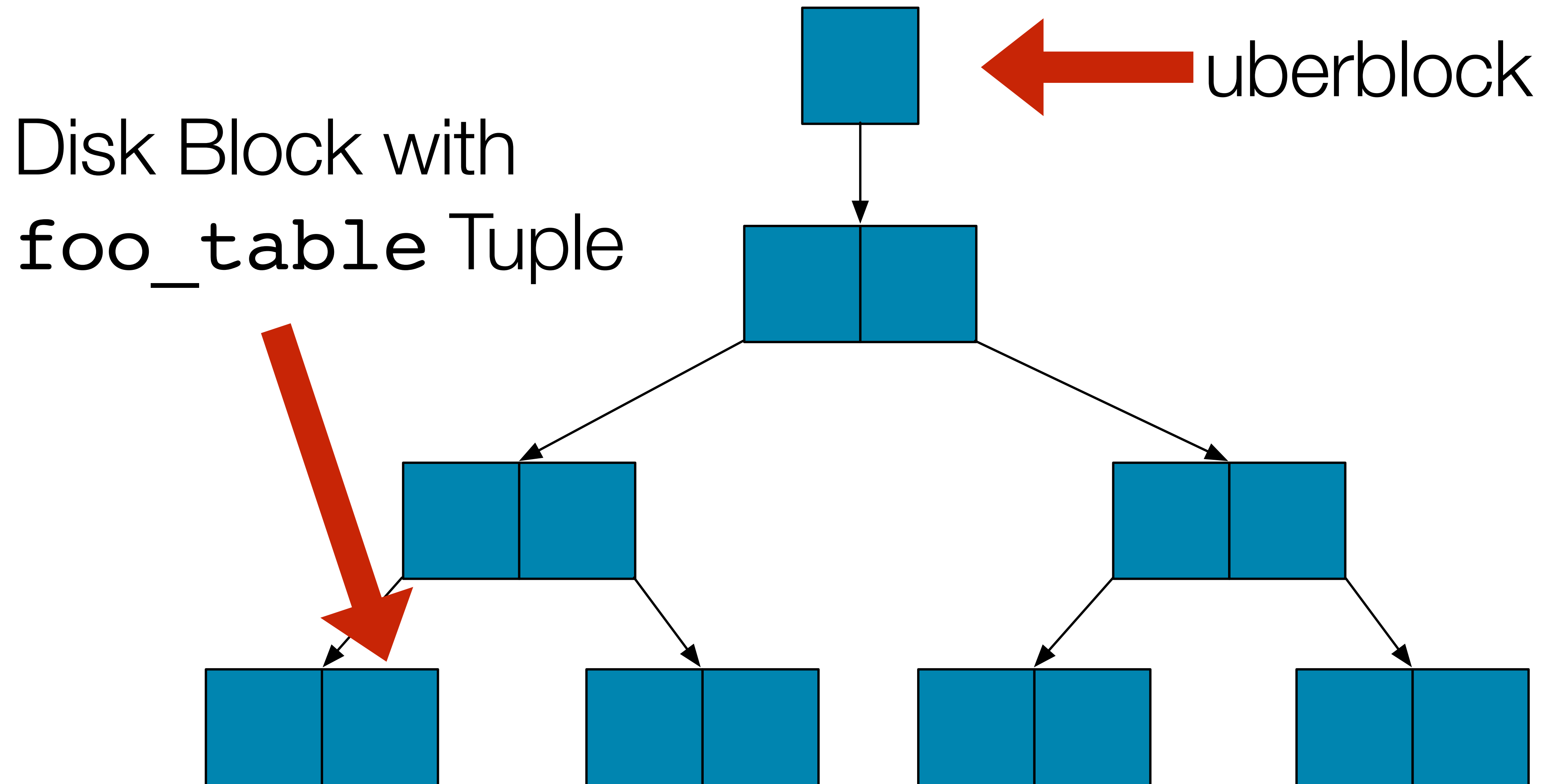
Log-Structured Filesystems: Top-Down

Log-Structured Filesystems: Top-Down

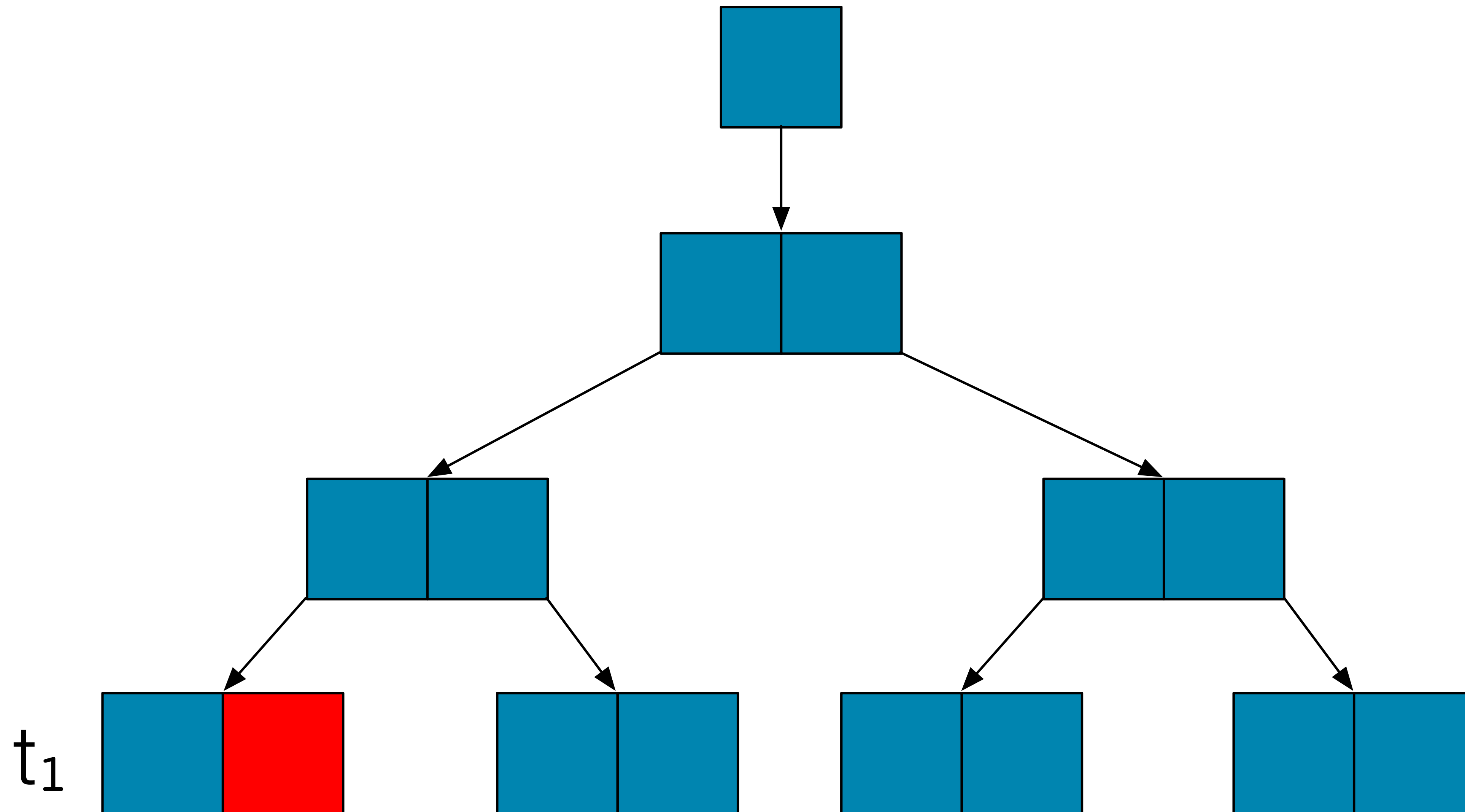
Disk Block with
`foo_table` Tuple



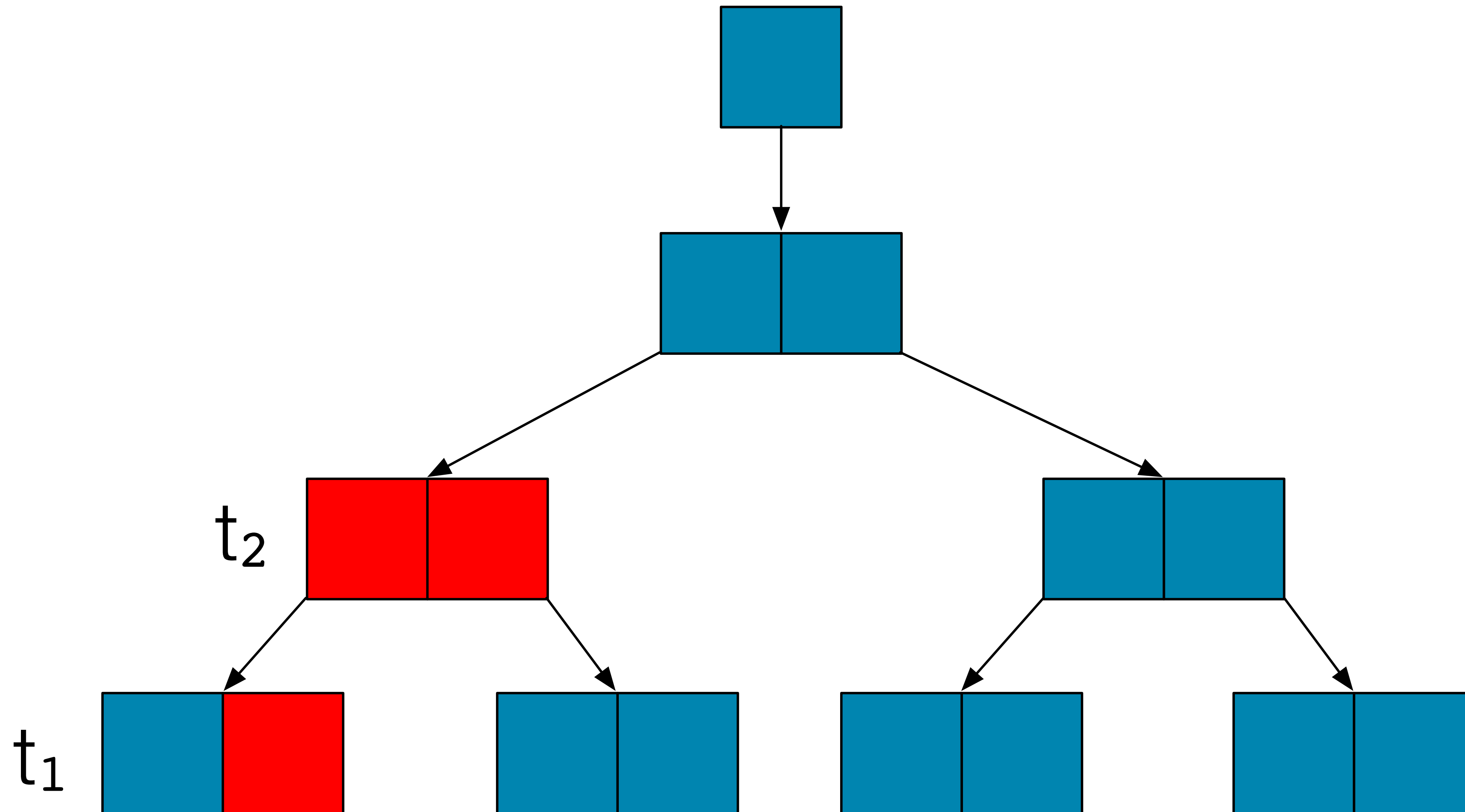
ZFS: User Data Block Lookup via ZFS Posix Layer



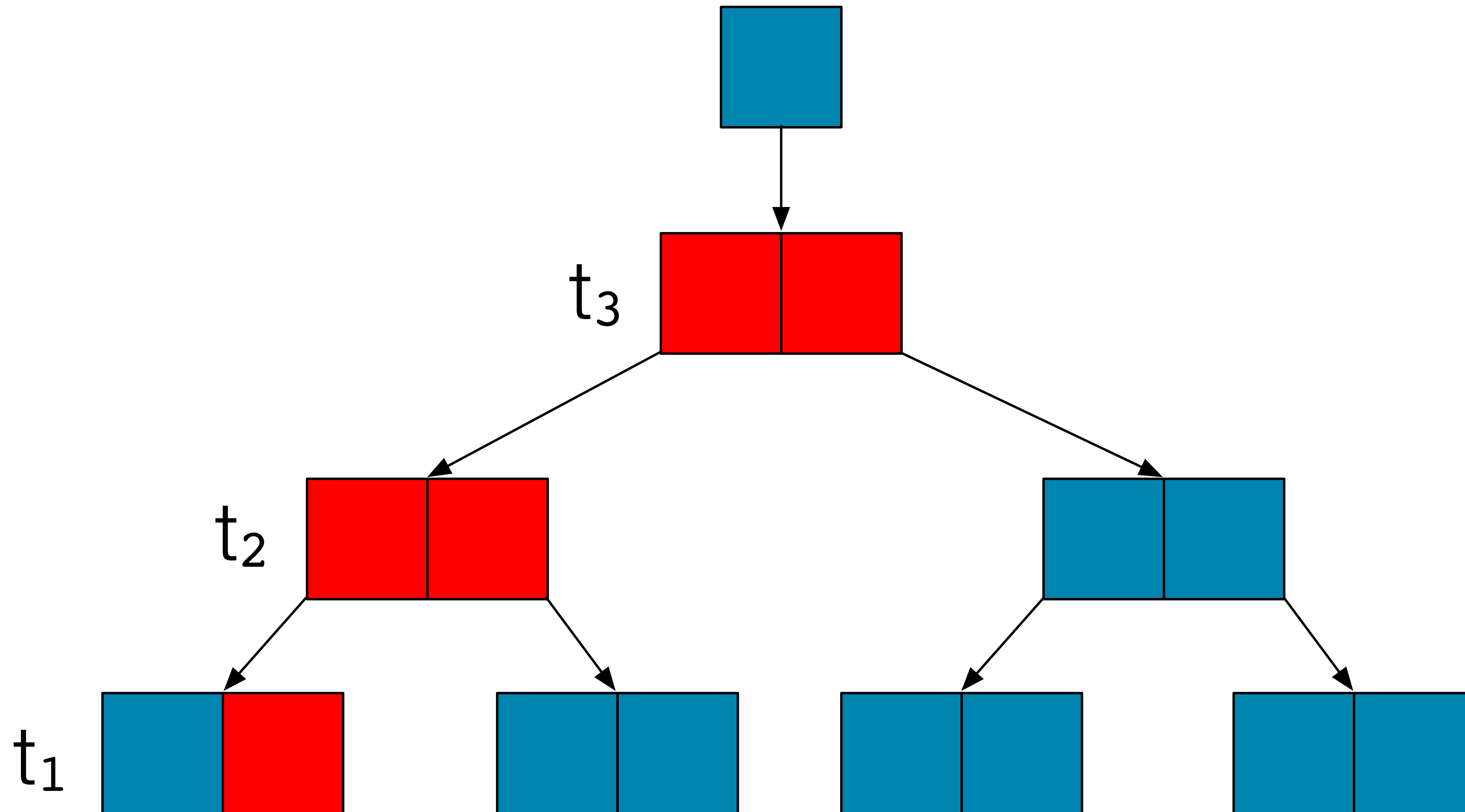
ZFS: User Data + File dnode



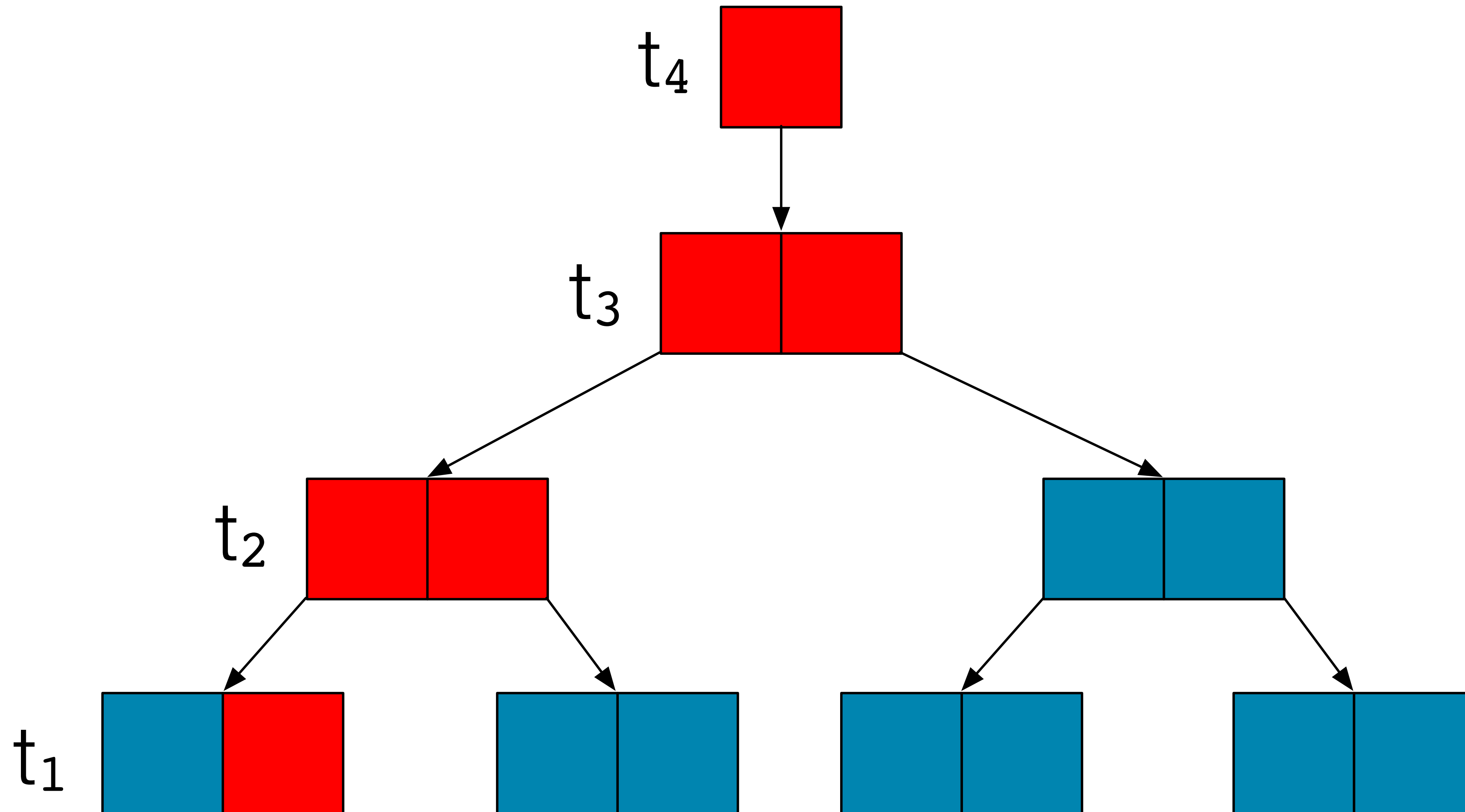
ZFS: Object Set



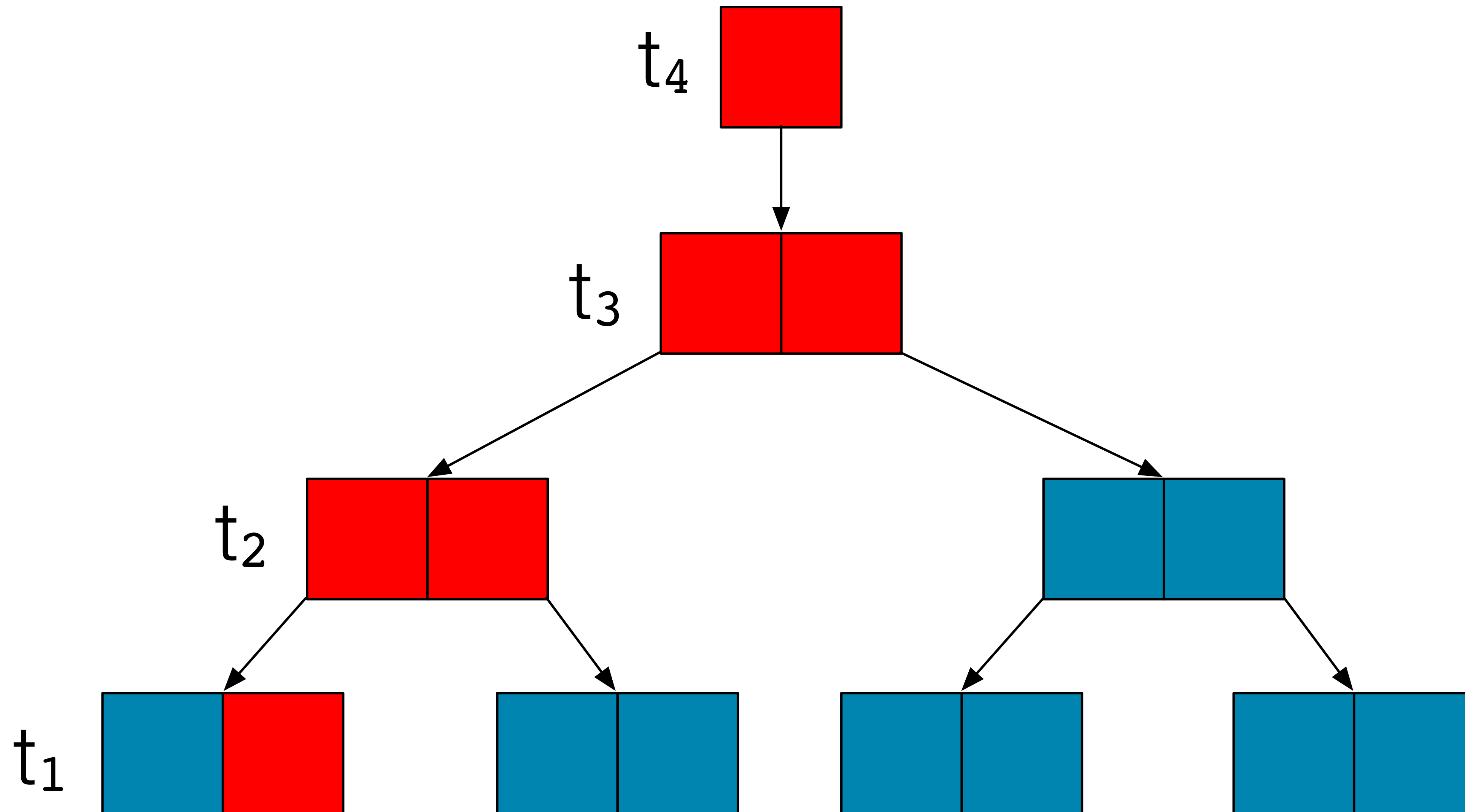
ZFS: Meta-Object Set Layer



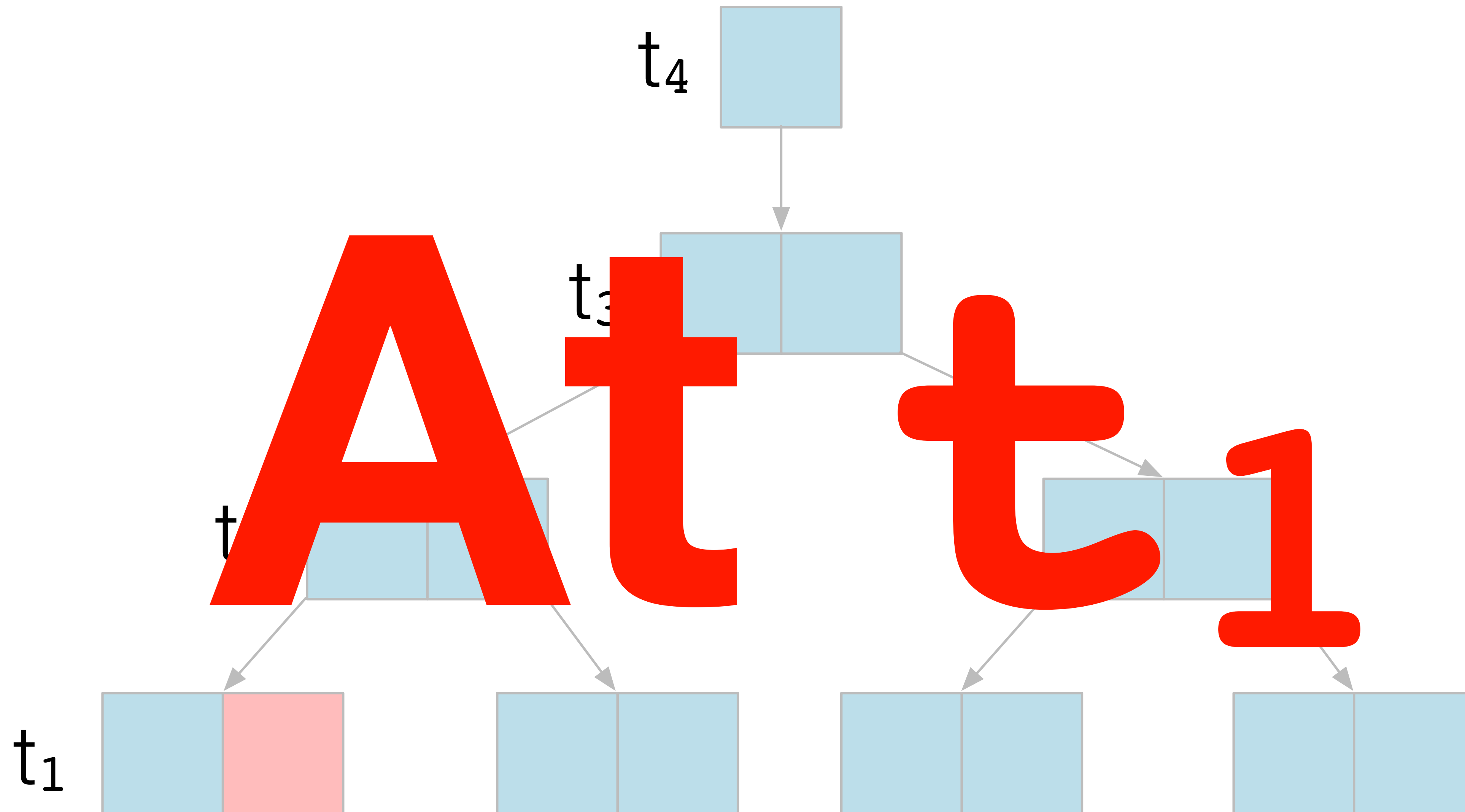
ZFS: Uberblock



At what point did the filesystem become inconsistent?



At what point could the filesystem become inconsistent?



How? I lied while explaining the situation. Alternate Truth.

Neglected to highlight **ZFS is Copy-On-Write** (read: knowingly committed perjury in front of a live audience)

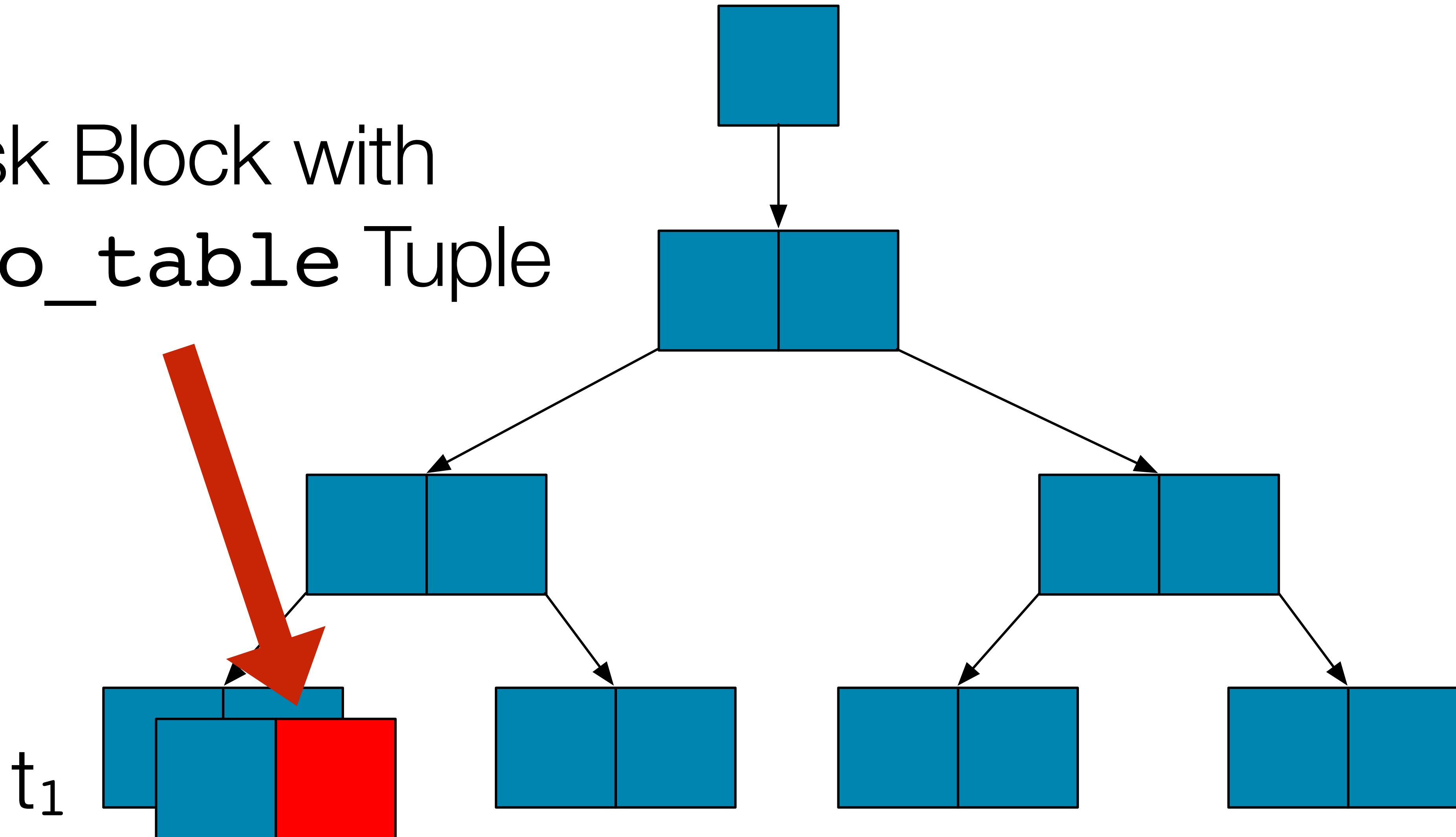
ZFS is Copy-On-Write

What what's not been deleted and on disk is immutable.

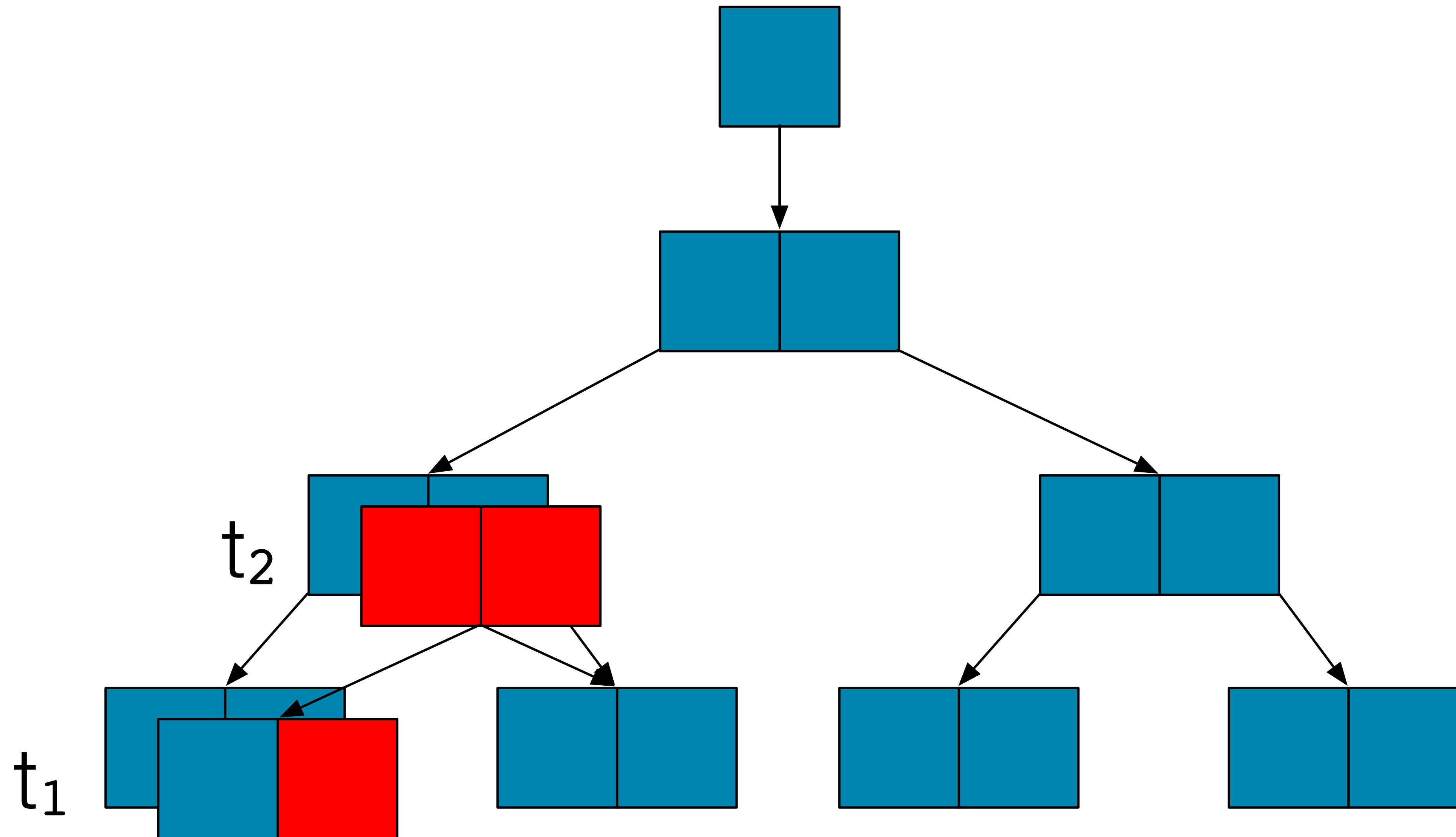
(read: I nearly committed perjury in front of a live audience by knowingly withholding vital information)

ZFS is Copy-On-Write

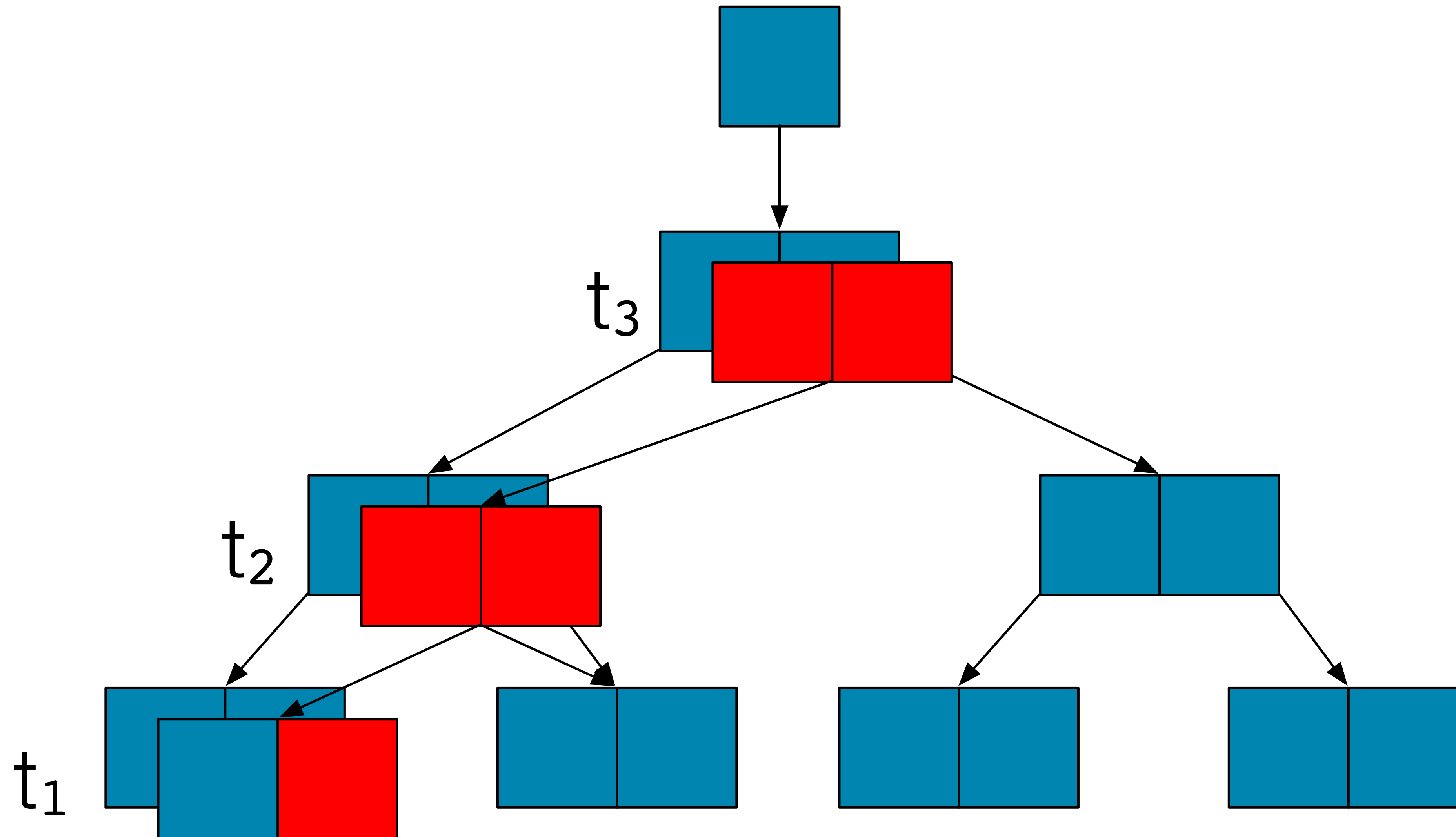
Disk Block with
`foo_table` Tuple



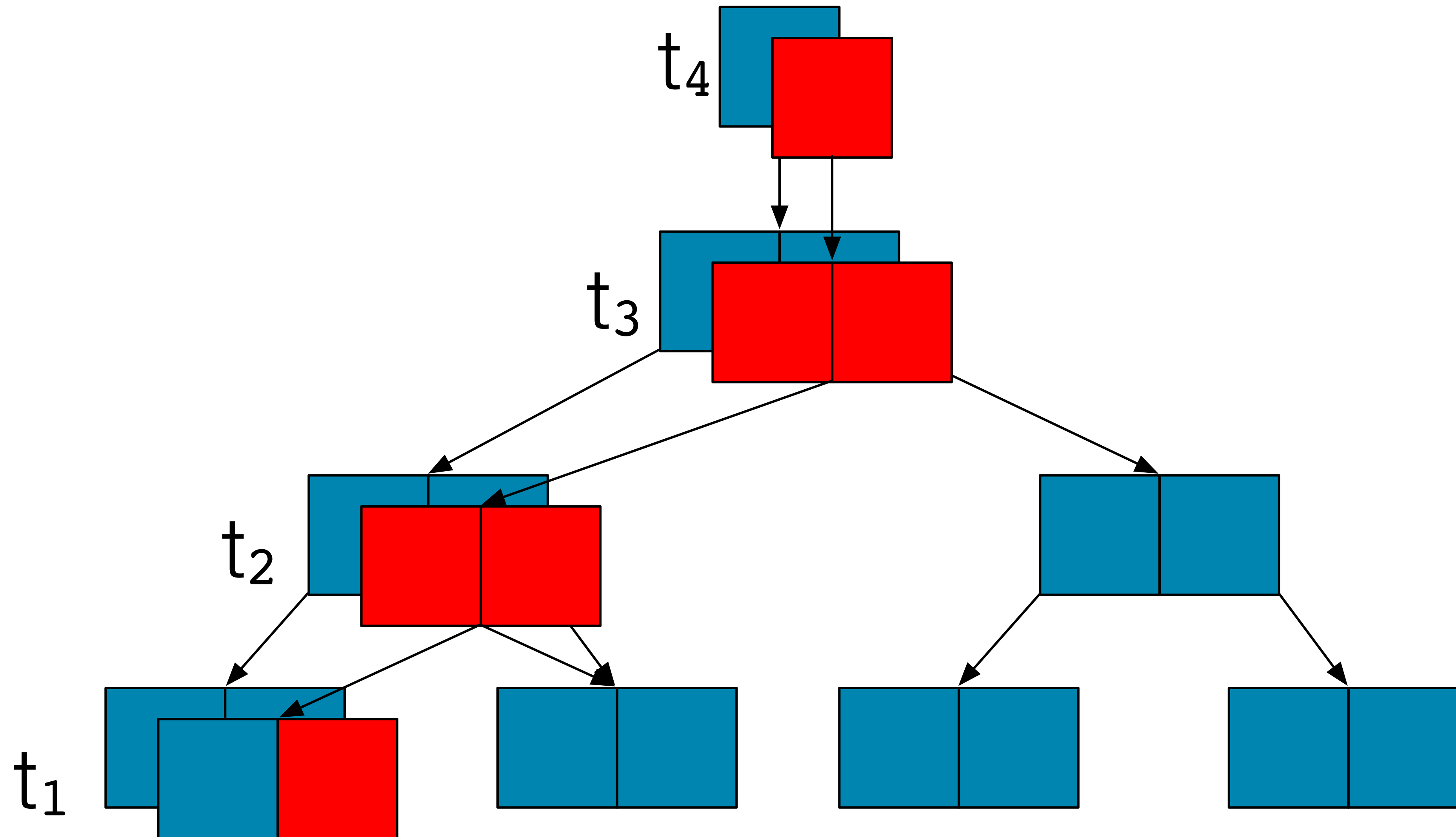
At what point did the filesystem become inconsistent?



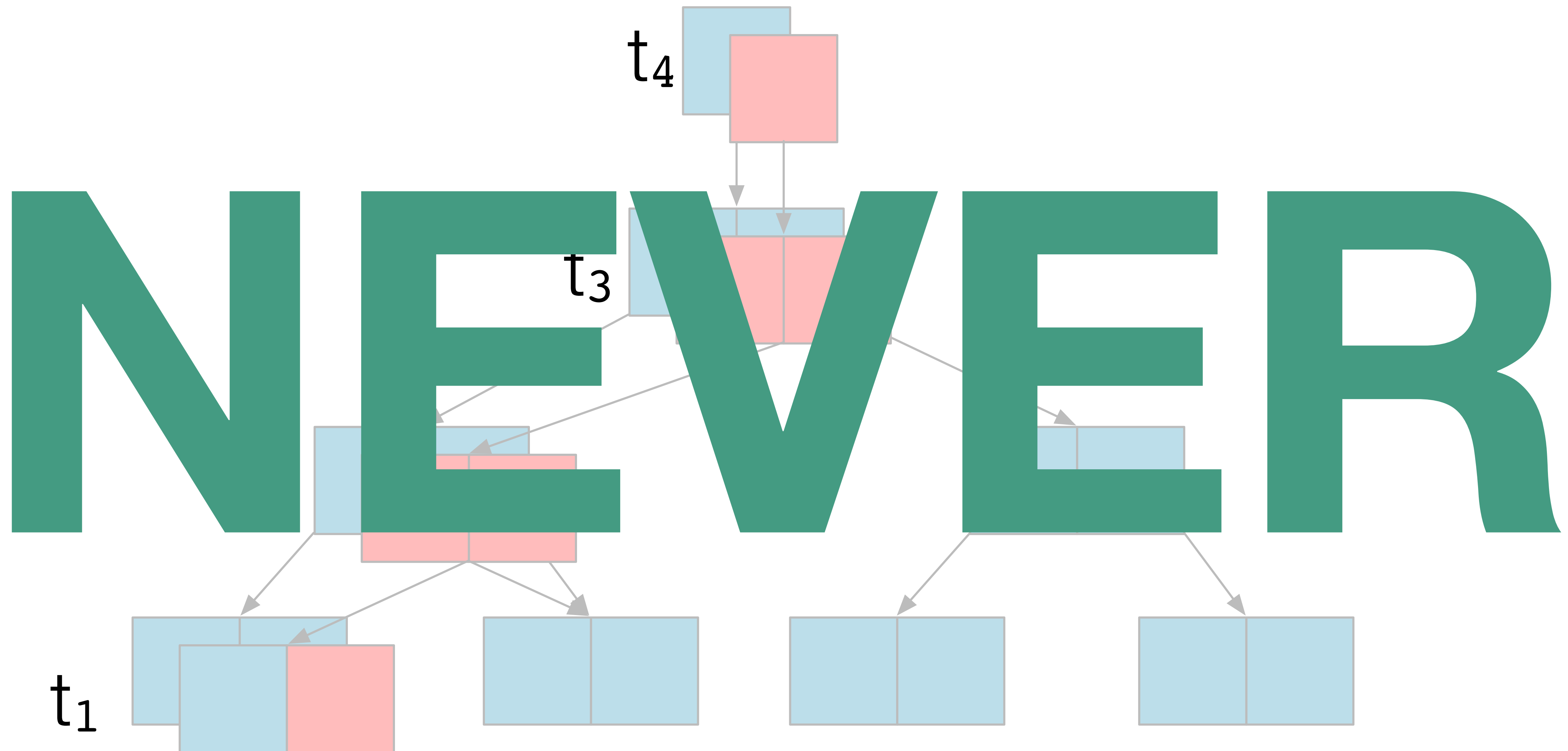
At what point did the filesystem become inconsistent?



At what point did the filesystem become inconsistent?



At what point could the filesystem become inconsistent?



TIL about ZFS: Transactions and Disk Pages

- Transaction groups are flushed to disk every N seconds (defaults to 5s)
- A transaction group (`txg`) in ZFS is called a "checkpoint"
- User Data can be modified as its written to disk
- All data is checksummed
- Compression should be enabled by default

ZFS Tip: ALWAYS enable compression

```
$ zfs get compression
NAME                PROPERTY            VALUE                SOURCE
rpool                compression          off                  default
rpool/root           compression          off                  default
$ sudo zfs set compression=lz4 rpool
$ zfs get compression
NAME                PROPERTY            VALUE                SOURCE
rpool                compression          lz4                  local
rpool/root           compression          lz4                  inherited from rpool
```

- Across ~7PB of PostgreSQL and mixed workloads and applications: compression ratio of ~2.8:1 was the average.
- Have seen >100:1 compression on some databases
(*cough* this data probably didn't belong in a database *cough*)
- Have seen as low as 1.01:1

ZFS Tip: ALWAYS enable compression

```
$ zfs get compression
NAME                PROPERTY            VALUE                SOURCE
rpool               compression          off                  default
rpool/root          compression          off                  default
$ sudo zfs set compression=lz4 rpool
$ zfs get compression
NAME                PROPERTY            VALUE                SOURCE
rpool               compression          lz4                  local
rpool/root          compression          lz4                  inherited from rpool
```

I have yet to see compression slow down benchmarking results or real world workloads. My experience is with:

- spinning rust (7.2K RPM, 10K RPM, and 15KRPM)
- fibre channel connected SANs
- SSDs
- NVME

ZFS Tip: ALWAYS enable compression

```
$ zfs get compressratio
```

NAME	PROPERTY	VALUE	SOURCE
rpool	compressratio	1.64x	—
rpool/db	compressratio	2.58x	—
rpool/db/pgdb1-10	compressratio	2.61x	—
rpool/root	compressratio	1.62x	—

- Use `lz4` by default everywhere.
- Use `gzip-9` only for archive servers
- Never mix-and-match compression where you can't suffer the consequences of lowest-common-denominator performance
- Anxious to see `zstandard` support (I'm looking at you Allan Jude)

ZFS Perk: Data Locality

- Data written at the same time is stored near each other because it's frequently part of the same record
- Data can now pre-fault into kernel cache (ZFS ARC) by virtue of the temporal adjacency of the related `pwrite(2)` calls
- Write locality + `compression=lz4` + `pg_repack` == PostgreSQL Dream Team

ZFS Perk: Data Locality

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- Write locality + `compression=lz4` + `pg_repack` == PostgreSQL Dream Team

If you don't know what `pg_repack` is, figure out how to move into a database environment that supports `pg_repack` and use it regularly.

https://reorg.github.io/pg_repack/ && https://github.com/reorg/pg_repack/

Extreme ZFS Warning: Purge all memory of dedup

- This is not just my recommendation, it's also from the community and author of the feature.
- These are not the droids you are looking for
- Do not pass Go
- Do not collect \$200
- Go straight to system unavailability jail
- The feature works, but you run the risk of bricking your ZFS server.

Ask after if you are curious, but here's a teaser:

What do you do if the `dedup` hash tables don't fit in RAM?

Bitrot is a Studied Phenomena

A Large-Scale Study of Flash Memory Failures in the Field

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ABSTRACT

Servers use flash memory based solid state drives (SSDs) as a high-performance alternative to hard disk drives to store persistent data. Unfortunately, recent increases in flash density have also brought about decreases in chip-level reliability. In a data center environment, flash-based SSD failures can lead to downtime and, in the worst case, data loss. As a result, it is important to understand flash memory reliability characteristics over flash lifetime in a realistic production data center environment running modern applications and system software.

This paper presents the first large-scale study of flash-based SSD reliability in the field. We analyze data collected across a majority of flash-based solid state drives at Facebook data centers over nearly four years and many millions of operational hours in order to understand failure properties and trends of flash-based SSDs. Our study considers a variety of SSD characteristics, including: the amount of data written to and read from flash chips; how data is mapped within the SSD address space; the amount of data copied, erased, and discarded by the flash controller; and flash board temperature and bus power.

Based on our field analysis of how flash memory errors manifest when running modern workloads on modern SSDs, this paper is the first to make several major observations: (1) SSD failure rates do *not* increase monotonically with *flash chip* wear; instead they go through several distinct periods corresponding to how failures emerge and are subsequently detected, (2) the effects of read disturbance errors are *not* prevalent in the field, (3) sparse logical data layout across an SSD's physical address space (e.g., non-contiguous data), as measured by the amount of metadata required to track logical address translations stored in an SSD-internal DRAM buffer, can greatly affect SSD failure rate, (4) higher temperatures lead to higher failure rates, but techniques that throttle SSD operation appear to greatly *reduce* the negative reliability impact of higher temperatures, and (5) data written by the operating system to flash-based SSDs does *not* always accurately

Categories and Subject Descriptors

B.3.4. [Hardware]: Memory Structures—*Reliability, Testing, and Fault-Tolerance*

Keywords

flash memory; reliability; warehouse-scale data centers

1. INTRODUCTION

Servers use flash memory for persistent storage due to the low access latency of flash chips compared to hard disk drives. Historically, flash capacity has lagged behind hard disk drive capacity, limiting the use of flash memory. In the past decade, however, advances in NAND flash memory technology have increased flash capacity by more than 1000 \times . This rapid increase in flash capacity has brought both an increase in flash memory use and a decrease in flash memory reliability. For example, the number of times that a cell can be reliably programmed and erased before wearing out and failing dropped from 10,000 times for 50 nm cells to only 2,000 times for 20 nm cells [28]. This trend is expected to continue for newer generations of flash memory. Therefore, if we want to improve the operational lifetime and reliability of flash memory-based devices, we must first fully understand their failure characteristics.

In the past, a large body of prior work examined the failure characteristics of flash cells in controlled environments using small numbers e.g., tens of raw flash chips (e.g., [36, 23, 21, 27, 22, 25, 16, 33, 14, 5, 18, 4, 24, 40, 41, 26, 31, 30, 37, 6, 11, 10, 7, 13, 9, 8, 12, 20]). This work quantified a variety of flash cell failure modes and formed the basis of the community's understanding of flash cell reliability. Yet prior work was limited in its analysis because these studies: (1) were conducted on small numbers of raw flash chips accessed in adversarial manners over short amounts of time, (2) did not examine failures when using real applications running on modern servers and

Bitrot is a Studied Phenomena

3.1 Bit Error Rate

The bit error rate (BER) of an SSD is the rate at which errors occur relative to the amount of information that is transmitted from/to the SSD. BER can be used to gauge the reliability of data transmission across a medium. We measure the uncorrectable bit error rate (UBER) from the SSD as:

$$UBER = \frac{\text{Uncorrectable errors}}{\text{Bits accessed}}$$

For flash-based SSDs, UBER is an important reliability metric that is related to the SSD lifetime. SSDs with high UBERs are expected to have more failed cells and encounter more (severe) errors that may potentially go undetected and corrupt

Bitrot is a Studied Phenomena

data than SSDs with low UBERs. Recent work by Grupp et al. examined the BER of *raw* MLC flash chips (without performing error correction in the flash controller) in a controlled environment [20]. They found the raw BER to vary from 1×10^{-1} for the least reliable flash chips down to 1×10^{-8} for the most reliable, with most chips having a BER in the 1×10^{-6} to 1×10^{-8} range. Their study did *not* analyze the effects of the use of chips *in SSDs under real workloads and system software*.

Table 1 shows the UBER of the platforms that we examine. We find that for flash-based SSDs used in servers, the UBER ranges from 2.6×10^{-9} to 5.1×10^{-11} . While we expect that the UBER of the SSDs that we measure (which correct small errors, perform wear leveling, and are not at the end of their rated life but still being used in a production environment) should be less than the raw chip BER in Grupp et al.’s study (which did not correct any errors in the flash controller, exercised flash chips until the end of their rated life, and accessed flash chips in an adversarial manner), we find that in some cases the BERs were within around one order of magnitude of each other. For example, the UBER of Platform B in our study, 2.6×10^{-9} , comes close to the lower end of the raw chip BER range reported in prior work, 1×10^{-8} .

Bitrot is a Studied Phenomena

Figure 2 (bottom) shows the average yearly uncorrectable error rate among SSDs within the different platforms – the sum of errors that occurred on all servers within a platform over 12 months ending in November 2014 divided by the number of servers in the platform. The yearly rates of uncorrectable errors on the SSDs we examined range from 15,128 for Platform D to 978,806 for Platform B. The older Platforms A and B have a higher error rate than the younger Platforms C through F, suggesting that the incidence of uncorrectable errors increases as SSDs are utilized more. We will examine this relationship further in Section 4.

Platform B has a much higher average yearly rate of uncorrectable errors (978,806) compared to the other platforms (the second highest amount, for Platform A, is 106,678). We find that this is due to a small number of SSDs having a much

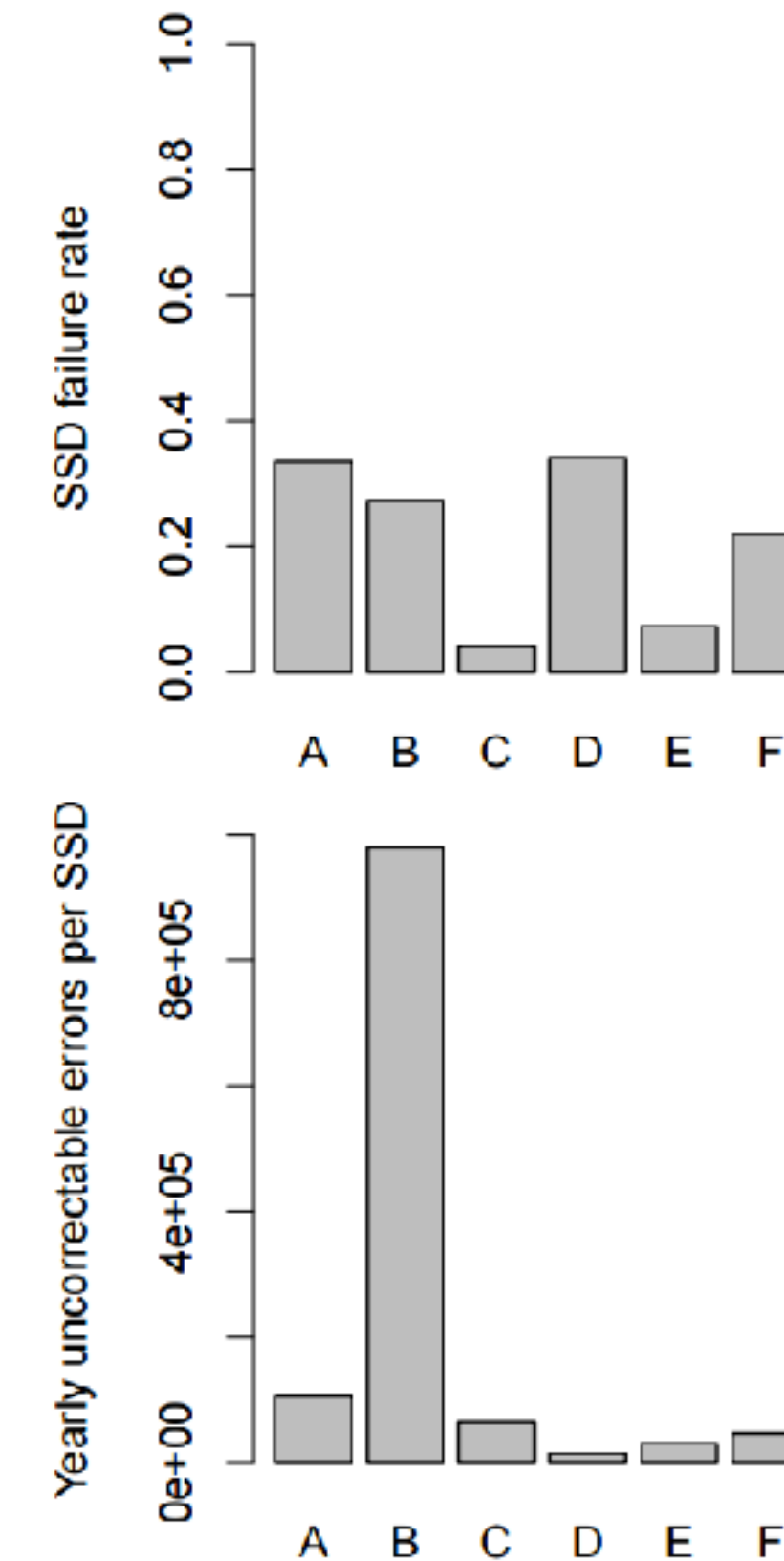


Figure 2: The failure rate (top) and average yearly uncorrectable errors per SSD (bottom), among SSDs within each platform.

TIL: Bitrot is here

- TL;DR: 4.2% -> 34% of SSDs have one UBER per year

$(1 - (1 - \text{uberRate})^{\text{numDisks}}) = \text{Probability of UBER/server/year}$

$(1 - (1 - 0.042)^{20}) = 58\%$

$(1 - (1 - 0.34)^{20}) = 99.975\%$

Highest quality SSD drives on the market



Lowest quality commercially viable SSD drives on the market

Causes of bitrot are Internal and External

External Factors for UBER on SSDs:

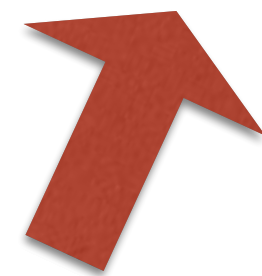
- Temperature
- Bus Power Consumption
- Data Written by the System Software
- Workload changes due to SSD failure

...except maybe they can.



Take Care of your bits

```
$ zpool status tank | head -n 3
pool: tank
state: ONLINE
scan: scrub repaired 4.50K in 53h44m with 0 errors on Tue May 26 21:36:26 2015
```



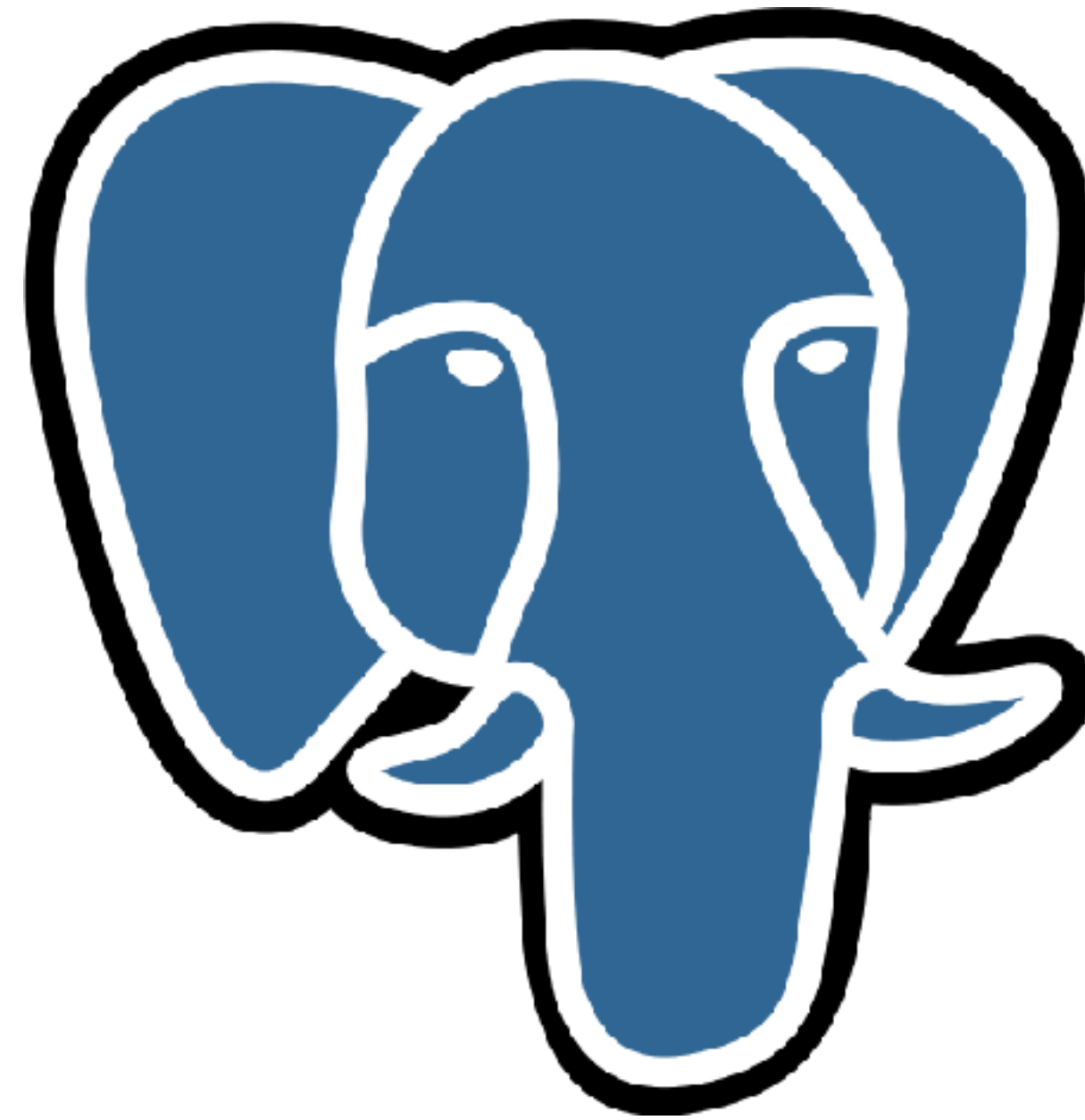
Answer their cry for help.

Take Care of your bits

Similar studies and research exist for:

- Fibre Channel
- SAS
- SATA
- Tape
- SANs
- Cloud Object Stores

So what about PostgreSQL?



"...I told you all of that, so I can tell you this..."

ZFS Terminology: VDEV

VDEV | vē-dēv

noun

a virtual device

- Physical drive redundancy is handled at the VDEV level
- Zero or more physical disks arranged like a RAID set:
 - mirror
 - stripe
 - raidz
 - raidz2
 - raidz3

ZFS Terminology: **zpool**

zpool | zē-pōol

noun

an abstraction of physical storage made up of a set of VDEVs

Loose a VDEV, loose the zpool.

ZFS Terminology: ZPL

ZPL | zē-pē-el

noun

ZFS POSIX Layer

- Layer that handles the impedance mismatch between POSIX filesystem semantics and the ZFS "object database."

ZFS Terminology: ZIL

ZIL | zil

noun

ZFS Intent Log

- The ZFS analog of PostgreSQL's WAL
- If you use a ZIL:
 - Use disks that have low-latency writes
 - Mirror your ZIL
- If you lose your ZIL, whatever data had not made it to the main data disks will be lost. The PostgreSQL equivalent of: `rm -rf pg_xlog/`

ZFS Terminology: ARC

ARC | ärk

noun

Adaptive Replacement Cache

- ZFS's page cache
- ARC will grow or shrink to match use up all of the available memory

TIP: Limit ARC's max size to a percentage of physical memory minus the `shared_buffer` cache for PostgreSQL minus the kernel's memory overhead.

ZFS Terminology: Datasets

dataset | dædəˌset

noun

A filesystem or volume ("zvol")

- A ZFS filesystem dataset uses the underlying **zpool**
- A dataset belongs to one and only one **zpool**
- Misc tunables, including compression and quotas are set on the dataset level

ZFS Terminology: The Missing Bits

ZAP	ZFS Attribute Processor
DMU	Data Management Unit
DSL	Dataset and Snapshot Layer
SPA	Storage Pool Allocator
ZVOL	ZFS Volume
ZIO	ZFS I/O
RAIDZ	RAID with variable-size stripes
L2ARC	Level 2 Adaptive Replacement Cache
record	unit of user data, think RAID stripe size


```
$ sudo zfs list
```

NAME	USED	AVAIL	REFER	MOUNTPOINT
rpool	818M	56.8G	96K	none
rpool/root	817M	56.8G	817M	/

```
$ ls -lA -d /db
```

```
ls: cannot access '/db': No such file or directory
```

```
$ sudo zfs create rpool/db -o mountpoint=/db
```

```
$ sudo zfs list
```

NAME	USED	AVAIL	REFER	MOUNTPOINT
rpool	818M	56.8G	96K	none
rpool/db	96K	56.8G	96K	/db
rpool/root	817M	56.8G	817M	/

```
$ ls -lA /db
```

```
total 9
```

```
drwxr-xr-x  2 root root  2 Mar  2 18:06 ./
```

```
drwxr-xr-x 22 root root 24 Mar  2 18:06 ../
```

Storage Management

- Running out of disk space is bad, m'kay?
- Block file systems reserve ~8% of the disk space above 100%
- At ~92% capacity the performance of block allocators change from "performance optimized" to "space optimized" (read: performance "drops").

Storage Management

- Running out of disk space is bad, m'kay?
- Block file systems reserve ~8% of the disk space above 100%
- At ~92% capacity the performance of block allocators change from "performance optimized" to "space optimized" (read: performance "drops").

ZFS doesn't have an artificial pool of free space: you have to manage that yourself.

```
$ sudo zpool list -H -o size  
59.6G  
$ sudo zpool list
```

The pool should never consume more than 80% of the available space

```
$ sudo zfs set quota=48G rpool/db
```

```
$ sudo zfs get quota rpool/db
```

NAME	PROPERTY	VALUE	SOURCE
rpool/db	quota	48G	local

```
$ sudo zfs list
```

NAME	USED	AVAIL	REFER	MOUNTPOINT
rpool	818M	56.8G	96K	none
rpool/db	96K	48.0G	96K	/db
rpool/root	817M	56.8G	817M	/

Dataset Tuning Tips

- Disable `atime`
- Enable `compression`
- Tune the `recordsize`
- Consider tweaking the `primarycache`

```
# zfs get atime,compression,primarycache,recordsize rpool/db
NAME          PROPERTY      VALUE          SOURCE
rpool/db      atime         on             inherited from rpool
rpool/db      compression   lz4            inherited from rpool
rpool/db      primarycache  all           default
rpool/db      recordsize    128K          default

# zfs set atime=off rpool/db
# zfs set compression=lz4 rpool/db
# zfs set recordsize=16K rpool/db
# zfs set primarycache=metadata rpool/db
# zfs get atime,compression,primarycache,recordsize rpool/db
NAME          PROPERTY      VALUE          SOURCE
rpool/db      atime         off            local
rpool/db      compression   lz4            local
rpool/db      primarycache  metadata       local
rpool/db      recordsize    16K           local
```

Discuss: recordsize=16K

- Pre-fault next page: useful for sequential scans
- With **compression=lz4**, reasonable to expect ~3-4x pages worth of data in a single ZFS record

Anecdotes and Recommendations:

- Performed better in most workloads vs ZFS's prefetch
- Disabling prefetch isn't necessary, tends to still be a net win
- Monitor arc cache usage

Discuss: **primarycache=metadata**

- metadata instructs ZFS's ARC to only cache metadata (e.g. **dnode** entries), not page data itself
- Default: cache all data

Two different recommendations based on benchmark workloads:

- Enable **primarycache=all** where working set exceeds RAM
- Enable **primarycache=metadata** where working set fits in RAM

Discuss: `primarycache=metadata`

- metadata instructs ZFS's ARC to only cache metadata (e.g. `dnode` entries), not page data itself
- Default: cache all data
- Double-caching happens

Two different recommendations based on benchmark workloads:

- Enable **`primarycache=all`** where working set exceeds RAM
- Enable **`primarycache=metadata`** where working set fits in RAM

Reasonable Default anecdote: Cap max ARC size ~15%-25% physical RAM + ~50% RAM `shared_buffers`

2-4µs/pwrite(2)!!

```
# dtrace -s vfs-io-postgres.d
Latencies (ns)
postgres Write
```

value	----- Distribution -----	count
1024		0
2048	@@	1325
4096	@@@@@@@	267
8192	@@	72
16384		0
32768		0
65536		19
131072		2
262144		0

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Performance Wins

# zpool iostat tank 1						
pool	alloc	capacity		operations		bandwidth
		free	read	write	read	write
-----	-----	-----	-----	-----	-----	-----
tank	958G	9.94T	0	210K	1022	330M
tank	958G	9.94T	1	207K	4.99K	326M
tank	958G	9.94T	32	30.5K	79.9K	46.9M
tank	958G	9.94T	22	9.62K	202K	15.9M
tank	958G	9.94T	15	10.2K	169K	16.5M
tank	958G	9.94T	36	10.5K	198K	14.9M
tank	958G	9.94T	6	10.8K	39.4K	17.4M
tank	958G	9.94T	12	189K	209K	298M
tank	958G	9.94T	1	210K	7.96K	340M
tank	958G	9.94T	10	218K	23.0K	355M
tank	958G	9.94T	2	224K	4.49K	359M
tank	958G	9.94T	6	228K	12.5K	367M
tank	958G	9.94T	7	140K	53.4K	225M
tank	958G	9.94T	9	26.9K	40.9K	44.0M
tank	958G	9.94T	0	9.43K	0	13.9M
tank	958G	9.94T	0	9.69K	0	16.3M
tank	958G	9.94T	1	74.0K	3.49K	120M
tank	958G	9.94T	6	226K	17.0K	366M
tank	958G	9.94T	0	225K	0	385M
tank	958G	9.94T	0	176K	0	515M
tank	958G	9.94T	0	84.7K	0	382M
tank	958G	9.94T	0	39.6K	0	163M

Performance Wins

```
# zpool iostat tank 1
```

pool	alloc	capacity		operations		bandwidth
		free	read	write	read	write
-----	-----	-----	-----	-----	-----	-----
tank	958G	9.94T	0	210K	1022	330M
tank	958G	9.94T	1	207K	4.99K	326M
tank	958G	9.94T	32	30.5K	79.9K	46.9M
tank	958G	9.94T	22	9.62K	202K	15.9M
tank	958G	9.94T	15	10.2K	169K	16.5M
tank	958G	9.94T	36	10.5K	198K	14.9M
tank	958G	9.94T	6	10.8K	39.4K	17.4M
tank	958G	9.94T	12	189K	209K	298M
tank	958G	9.94T	1	210K	7.96K	340M
tank	958G	9.94T	10	218K	23.0K	355M
tank	958G	9.94T	2	224K	4.49K	359M
tank	958G	9.94T	6	228K	12.5K	367M
tank	958G	9.94T	7	140K	53.4K	225M
tank	958G	9.94T	9	26.9K	40.9K	44.0M
tank	958G	9.94T	0	9.43K	0	13.9M
tank	958G	9.94T	0	9.69K	0	16.3M
tank	958G	9.94T	1	74.0K	3.49K	120M
tank	958G	9.94T	6	226K	17.0K	366M
tank	958G	9.94T	0	225K	0	385M
tank	958G	9.94T	0	176K	0	515M
tank	958G	9.94T	0	84.7K	0	382M
tank	958G	9.94T	0	39.6K	0	163M

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Performance Wins

# zpool iostat tank 1						
		capacity		operations		bandwidth
pool	alloc	free	read	write	read	write
-----	-----	-----	-----	-----	-----	-----
tank	958G	9.94T	0	210K	1022	330M
tank	958G	9.94T	1	207K	4.99K	326M
tank	958G	9.94T	32	30.5K	79.9K	46.9M
tank	958G	9.94T	32	0.62K	202K	15.9M
tank	958G	9.94T	36	10.5K	198K	14.9M
tank	958G	9.94T	6	10.8K	39.4K	17.4M
tank	958G	9.94T	12	189K	209K	298M
tank	958G	9.94T	1	210K	7.96K	340M
tank	958G	9.94T	10	218K	23.0K	355M
tank	958G	9.94T	2	224K	4.49K	359M
tank	958G	9.94T	6	228K	12.5K	367M
tank	958G	9.94T	7	140K	53.4K	225M
tank	958G	9.94T	9	26.9K	40.9K	44.0M
tank	958G	9.94T	0	9.43K	0	13.9M
tank	958G	9.94T	0	9.69K	0	16.3M
tank	958G	9.94T	1	74.0K	3.49K	120M
tank	958G	9.94T	6	226K	17.0K	366M
tank	958G	9.94T	0	225K	0	385M
tank	958G	9.94T	0	176K	0	515M
tank	958G	9.94T	0	84.7K	0	382M
tank	958G	9.94T	0	39.6K	0	163M

P.S. This was observed on 10K RPM spinning rust.

ZFS Always has your back

- ZFS will checksum every read from disk
- A failed checksum will result in a fault and automatic data reconstruction
- Scrubs do background check of every record
- Schedule periodic scrubs
 - Frequently for new and old devices
 - Infrequently for devices in service between 6mo and 2.5yr

PSA: The "Compressed ARC" feature was added to catch checksum errors in **RAM**

Checksum errors are an early indicator of failing disks

Schedule Periodic Scrubs

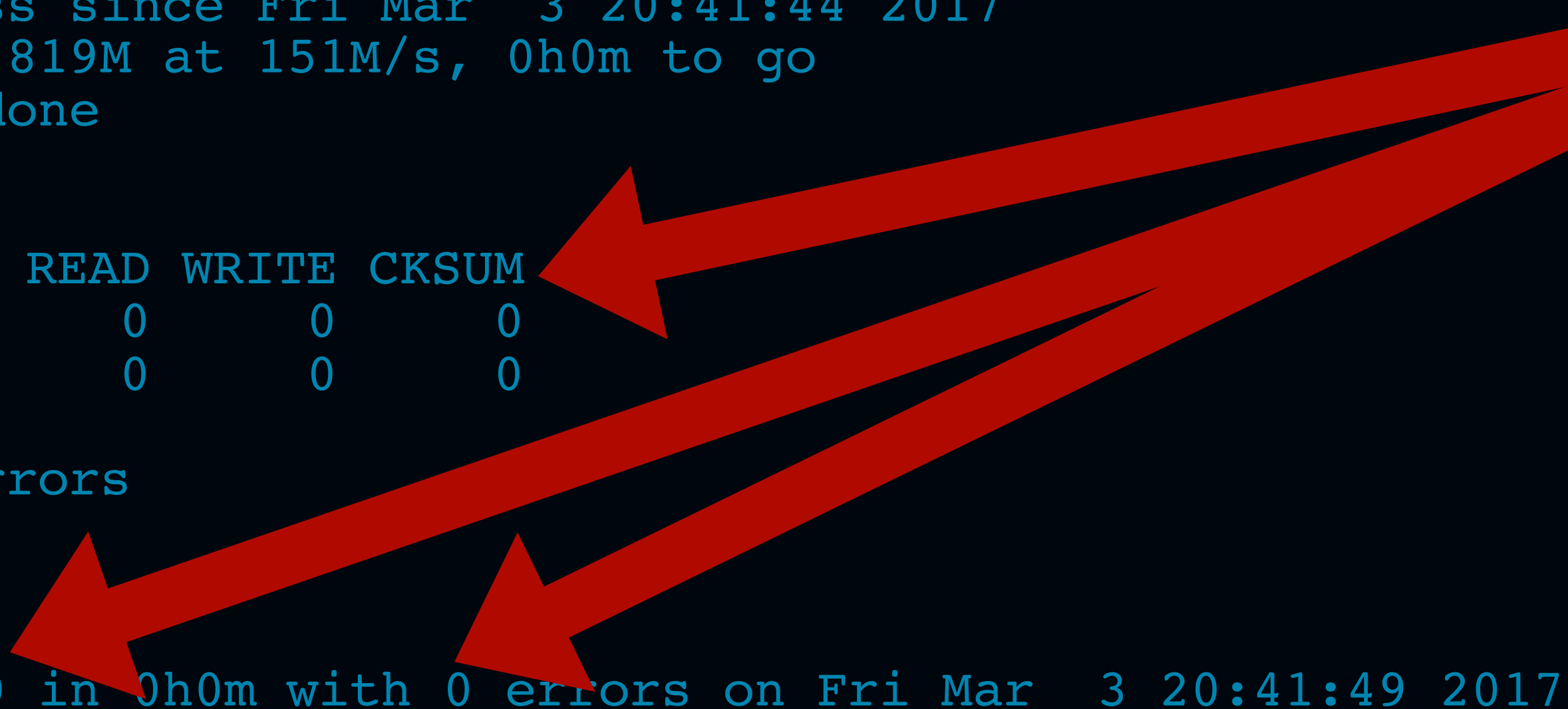
```
# zpool status
pool: rpool
state: ONLINE
scan: none requested
config:

    NAME      STATE      READ  WRITE CKSUM
    rpool     ONLINE      0     0     0
    sda1      ONLINE      0     0     0

errors: No known data errors
# zpool scrub rpool
# zpool status
pool: rpool
state: ONLINE
scan: scrub in progress since Fri Mar  3 20:41:44 2017
      753M scanned out of 819M at 151M/s, 0h0m to go
      0 repaired, 91.97% done
config:

    NAME      STATE      READ  WRITE CKSUM
    rpool     ONLINE      0     0     0
    sda1      ONLINE      0     0     0

errors: No known data errors
# zpool status
pool: rpool
state: ONLINE
scan: scrub repaired 0 in 0h0m with 0 errors on Fri Mar  3 20:41:49 2017
```



Non-zero on
any of these
values is bad™

One dataset per database

- Create one ZFS dataset per database instance
- General rules of thumb:
 - Use the same dataset for **\$PGDATA/** and **pg_xlogs/**
 - Set a reasonable quota
 - Optional: reserve space to guarantee minimal available space

Checksum errors are an early indicator of failing disks

One dataset per database

```
# zfs list
NAME                USED    AVAIL    REFER    MOUNTPOINT
rpool               819M    56.8G    96K      none
rpool/db            160K    48.0G    96K      /db
rpool/root          818M    56.8G    818M     /

# zfs create rpool/db/pgdb1
# chown postgres:postgres /db/pgdb1
# zfs list
NAME                USED    AVAIL    REFER    MOUNTPOINT
rpool               819M    56.8G    96K      none
rpool/db            256K    48.0G    96K      /db
rpool/db/pgdb1       96K     48.0G    96K      /db/pgdb1
rpool/root          818M    56.8G    818M     /

# zfs set reservation=1G rpool/db/pgdb1
# zfs list
NAME                USED    AVAIL    REFER    MOUNTPOINT
rpool               1.80G    55.8G    96K      none
rpool/db            1.00G    47.0G    96K      /db
rpool/db/pgdb1       96K     48.0G    12.0M    /db/pgdb1
rpool/root          818M    55.8G    818M     /
```

initdb like a boss

```
# su postgres -c 'initdb --no-locale -E=UTF8 -n -N -D /db/pgdb1'
Running in noclean mode.  Mistakes will not be cleaned up.
The files belonging to this database system will be owned by user "postgres".
This user must also own the server process.

The database cluster will be initialized with locale "C".
The default text search configuration will be set to "english".


Data page checksums are disabled.

fixing permissions on existing directory /db/pgdb1 ... ok
creating subdirectories ... ok
```

- Encode using UTF8, sort using C
- Only enable locale when you know you need it
 - ~2x perf bump by avoiding calls to `iconv(3)` to figure out sort order
 - **DO NOT** use PostgreSQL checksums or compression

```
# zfs list -t snapshot
no datasets available
# pwd
/db/pgdb1
# find . | wc -l
895
# head -1 postmaster.pid
25114
# zfs snapshot rpool/db/pgdb1@pre-rm
# zfs list -t snapshot
NAME                                USED    AVAIL    REFER    MOUNTPOINT
rpool/db/pgdb1@pre-rm              0        -    12.0M    -
# psql -U postgres
psql (9.6.2)
Type "help" for help.

postgres=# \q
# rm -rf *
# ls -l | wc -l
0
# psql -U postgres
psql: FATAL:  could not open relation mapping file "global/pg_filenode.map":
No such file or directory
```



Guilty Pleasure
During Demos

Backups: Has Them

```
$ psql
psql: FATAL:  could not open relation mapping file "global/pg_filenode.map": No such file or directory
# cat postgres.log
LOG:  database system was shut down at 2017-03-03 21:08:05 UTC
LOG:  MultiXact member wraparound protections are now enabled
LOG:  database system is ready to accept connections
LOG:  autovacuum launcher started
FATAL:  could not open relation mapping file "global/pg_filenode.map": No such file or directory
LOG:  could not open temporary statistics file "pg_stat_tmp/global.tmp": No such file or directory
LOG:  could not open temporary statistics file "pg_stat_tmp/global.tmp": No such file or directory
...
LOG:  could not open temporary statistics file "pg_stat_tmp/global.tmp": No such file or directory
LOG:  could not open file "postmaster.pid": No such file or directory
LOG:  performing immediate shutdown because data directory lock file is invalid
LOG:  received immediate shutdown request
LOG:  could not open temporary statistics file "pg_stat/global.tmp": No such file or directory
WARNING:  terminating connection because of crash of another server process
DETAIL:  The postmaster has commanded this server process to roll back the current transaction and exit,
because another server process exited abnormally and possibly corrupted shared memory.
HINT:  In a moment you should be able to reconnect to the database and repeat your command.
LOG:  database system is shut down
# ll
total 1
drwx----- 2 postgres postgres 2 Mar  3 21:40 ./
drwxr-xr-x 3 root      root      3 Mar  3 21:03 ../
```

Restores: As Important as Backups

```
# zfs list -t snapshot
NAME                                USED  AVAIL  REFER  MOUNTPOINT
rpool/db/pgdb1@pre-rm             12.0M      -   12.0M      -
# zfs rollback rpool/db/pgdb1@pre-rm
# su postgres -c '/usr/lib/postgresql/9.6/bin/postgres -D /db/pgdb1'
LOG:  database system was interrupted; last known up at 2017-03-03 21:50:57 UTC
LOG:  database system was not properly shut down; automatic recovery in progress
LOG:  redo starts at 0/14EE7B8
LOG:  invalid record length at 0/1504150: wanted 24, got 0
LOG:  redo done at 0/1504128
LOG:  last completed transaction was at log time 2017-03-03 21:51:15.340442+00
LOG:  MultiXact member wraparound protections are now enabled
LOG:  database system is ready to accept connections
LOG:  autovacuum launcher started
```

Works all the time, every time, even with `kill -9`

(possible dataloss from ungraceful shutdown and IPC cleanup notwithstanding)

Clone: Test and Upgrade with Impunity

```
# zfs clone rpool/db/pgdb1@pre-rm rpool/db/pgdb1-upgrade-test
# zfs list -r rpool/db
```

NAME	USED	AVAIL	REFER	MOUNTPOINT
rpool/db	1.00G	47.0G	96K	/db
rpool/db/pgdb1	15.6M	48.0G	15.1M	/db/pgdb1
rpool/db/pgdb1-upgrade-test	8K	47.0G	15.2M	/db/pgdb1-upgrade-test

```
# echo "Test pg_upgrade"
# zfs destroy rpool/db/pgdb1-clone
# zfs clone rpool/db/pgdb1@pre-rm rpool/db/pgdb1-10
# echo "Run pg_upgrade for real"
# zfs promote rpool/db/pgdb1-10
# zfs destroy rpool/db/pgdb1
```

Works all the time, every time, even with `kill -9`

(possible dataloss from ungraceful shutdown and IPC cleanup notwithstanding)

100 Tip: Naming Conventions

- Use a short prefix not on the root filesystem (e.g. /db)
- Encode the PostgreSQL major version into the dataset name
- Give each PostgreSQL cluster its own dataset (e.g. pgdb01)
- Optional but recommended:
 - one database per cluster
 - one app per database
 - encode environment into DB name
 - encode environment into DB username

Suboptimal	Good
rpool/db/pgdb1	rpool/db/prod-db01-pg94
rpool/db/myapp-shard1	rpool/db/prod-myapp-shard001-pg95
rpool/db/dbN	rpool/db/prod-dbN-pg10

Be explicit: codify the tight coupling between PostgreSQL versions and `$PGDATA/`.

Defy Gravity

- Take and send snapshots to remote servers
- **zfs send** emits a snapshot to stdout: treat as a file or stream
- **zfs receive** reads a snapshot from stdin
- TIP: If available:
 - Use the **-s** argument to **zfs receive**
 - Use **zfs get receive_resume_token** on the receiving end to get the required token to resume an interrupted send: **zfs send -t <token>**

Unlimited flexibility. Compress, encrypt, checksum, and offsite to your heart's content.

```
# zfs send -v -L -p -e rpool/db/pgdb1@pre-rm > /dev/null
send from @ to rpool/db/pgdb1-10@pre-rm estimated size is 36.8M
total estimated size is 36.8M
TIME          SENT      SNAPSHOT
# zfs send -v -L -p -e \
    rpool/db/pgdb1-10@pre-rm | \
    zfs receive -v \
    rpool/db/pgdb1-10-receive
send from @ to rpool/db/pgdb1-10@pre-rm estimated size is 36.8M
total estimated size is 36.8M
TIME          SENT      SNAPSHOT
received 33.8MB stream in 1 seconds (33.8MB/sec)
# zfs list -t snapshot
```

NAME	USED	AVAIL	REFER
MOUNTPOINT			
rpool/db/pgdb1-10@pre-rm	8K	-	15.2M -
rpool/db/pgdb1-10-receive@pre-rm	0	-	15.2M -

Defy Gravity: Incrementally

- Use a predictable snapshot naming scheme
- Send snapshots incrementally
- Clean up old snapshots
- Use a monotonic snapshot number (a.k.a. "vector clock")

Remember to remove old snapshots.
Distributed systems bingo!

Defy Gravity: Incremental

```
# echo "Change PostgreSQL's data"
# zfs snapshot rpool/db/pgdb1-10@example-incremental-001
# zfs send -v -L -p -e \
    -i rpool/db/pgdb1-10@pre-rm \
    rpool/db/pgdb1-10@example-incremental-001 \
    > /dev/null
send from @pre-rm to rpool/db/pgdb1-10@example-incremental-001
estimated size is 2K
total estimated size is 2K
# zfs send -v -L -p -e \
    -i rpool/db/pgdb1-10@pre-rm \
    rpool/db/pgdb1-10@example-incremental-001 | \
    zfs receive -v \
    rpool/db/pgdb1-10-receive
send from @pre-rm to rpool/db/pgdb1-10@example-incremental-001
estimated size is 2K
total estimated size is 2K
receiving incremental stream of rpool/db/pgdb1-10@example-
incremental-001 into rpool/db/pgdb1-10-receive@example-incremental-001
received 312B stream in 1 seconds (312B/sec)
```

Defy Gravity: Vector Clock

```
# echo "Change more PostgreSQL's data: VACUUM FULL FREEZE"
# zfs snapshot rpool/db/pgdb1-10@example-incremental-002
# zfs send -v -L -p -e \
    -i rpool/db/pgdb1-10@example-incremental-001 \
    rpool/db/pgdb1-10@example-incremental-002 \
    > /dev/null
send from @example-incremental-001 to rpool/db/pgdb1-10@example-
incremental-002 estimated size is 7.60M
total estimated size is 7.60M
TIME          SENT      SNAPSHOT
# zfs send -v -L -p -e \
    -i rpool/db/pgdb1-10@example-incremental-001 \
    rpool/db/pgdb1-10@example-incremental-002 | \
    zfs receive -v \
    rpool/db/pgdb1-10-receive
send from @example-incremental-001 to rpool/db/pgdb1-10@example-
incremental-002 estimated size is 7.60M
total estimated size is 7.60M
receiving incremental stream of rpool/db/pgdb1-10@example-incremental-002
into rpool/db/pgdb1-10-receive@example-incremental-002
TIME          SENT      SNAPSHOT
received 7.52MB stream in 1 seconds (7.52MB/sec)
```

Defy Gravity: Cleanup

```
# zfs list -t snapshot -o name,used,refer
NAME                                     USED    REFER
rpool/db/pgdb1-10@example-incremental-001      8K    15.2M
rpool/db/pgdb1-10@example-incremental-002    848K    15.1M
rpool/db/pgdb1-10-receive@pre-rm              8K    15.2M
rpool/db/pgdb1-10-receive@example-incremental-001  8K    15.2M
rpool/db/pgdb1-10-receive@example-incremental-002   0    15.1M
# zfs destroy rpool/db/pgdb1-10-receive@pre-rm
# zfs destroy rpool/db/pgdb1-10@example-incremental-001
# zfs destroy rpool/db/pgdb1-10-receive@example-incremental-001
# zfs list -t snapshot -o name,used,refer
NAME                                     USED    REFER
rpool/db/pgdb1-10@example-incremental-002    848K    15.1M
rpool/db/pgdb1-10-receive@example-incremental-002   0    15.1M
```


Controversial: `logbias=throughput`

- Measure tps/qps
- Time duration of an outage (OS restart plus WAL replay, e.g. 10-20min)
- Measure cost of back pressure from the DB to the rest of the application
- Use a **`txg`** timeout of 1 second

Position: since ZFS will never be inconsistent and therefore PostgreSQL will never lose integrity, 1s of actual data loss is a worthwhile tradeoff for a ~10x performance boost in write-heavy applications.

Rationale: loss aversion costs organizations more than potentially losing 1s of data. Back pressure is a constant cost the rest of the application needs to absorb due to continual **`fsync(2)`**'ing of WAL data. Architectural cost and premature engineering costs need to be factored in. Penny-wise, pound foolish.

Controversial: logbias=throughput

```
# cat /sys/module/zfs/parameters/zfs_txg_timeout
5
# echo 1 > /sys/module/zfs/parameters/zfs_txg_timeout
# echo 'options zfs zfs_txg_timeout=1' >> /etc/modprobe.d/zfs.conf
# psql -c 'ALTER SYSTEM SET synchronous_commit=off'
ALTER SYSTEM
# zfs set logbias=throughput rpool/db
```

Thank You

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 - Sergio Murilo
 - Chris Schneider
 - Filippos Kalamidas
 - Jose Finotto

QUESTIONS?

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